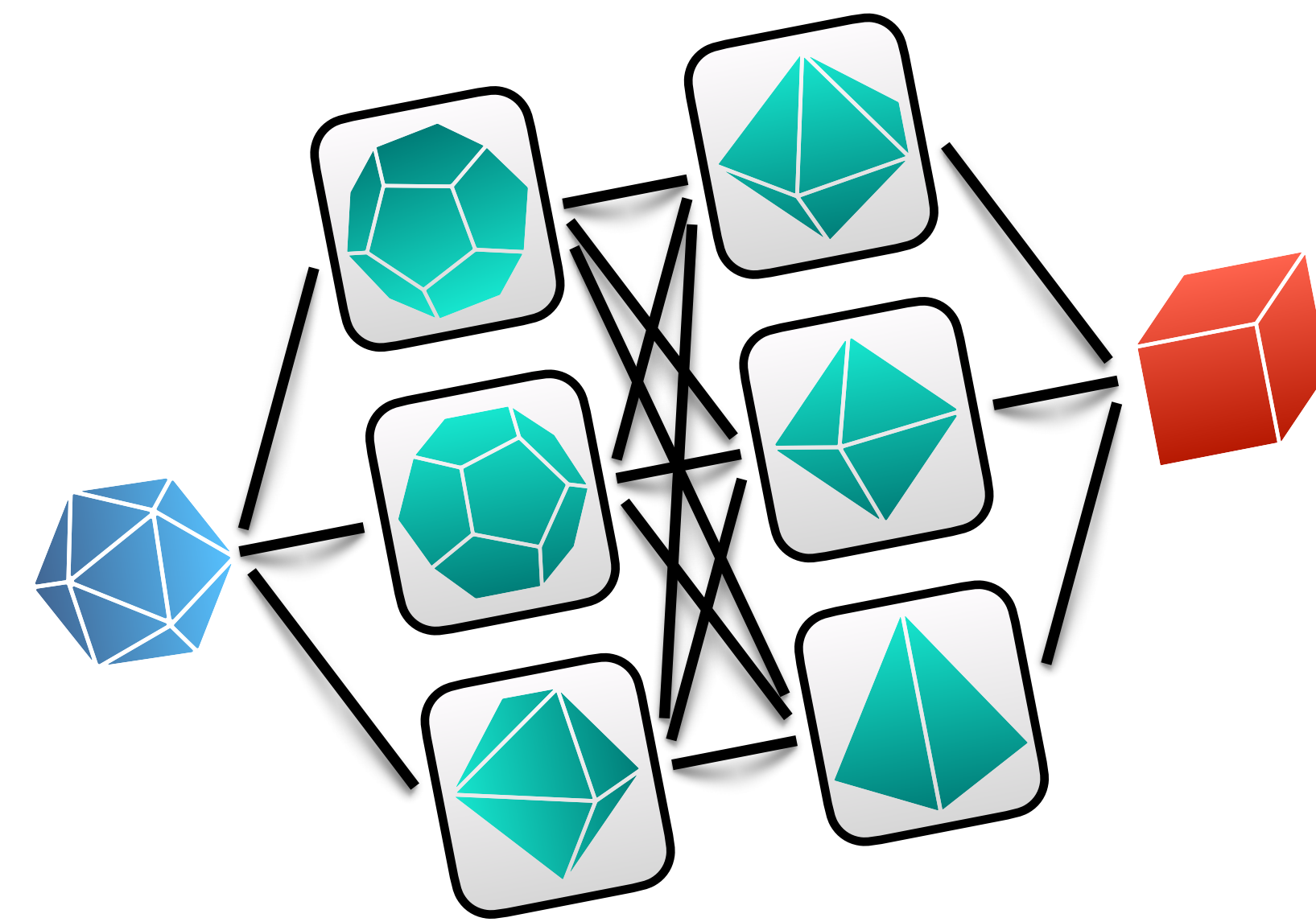




Did I forget to hit
record? Please
remind me!



Algorithm Configuration and friends

MIE1666: Machine Learning for Mathematical Optimization

Largely based on Hoos, Holger H. "Automated algorithm configuration and parameter tuning." Autonomous search. Springer, Berlin, Heidelberg, 2011. 37-71.

Some examples from Chapter 13 of Integer Programming by Wolsey

Elias B. Khalil — 04/10/21



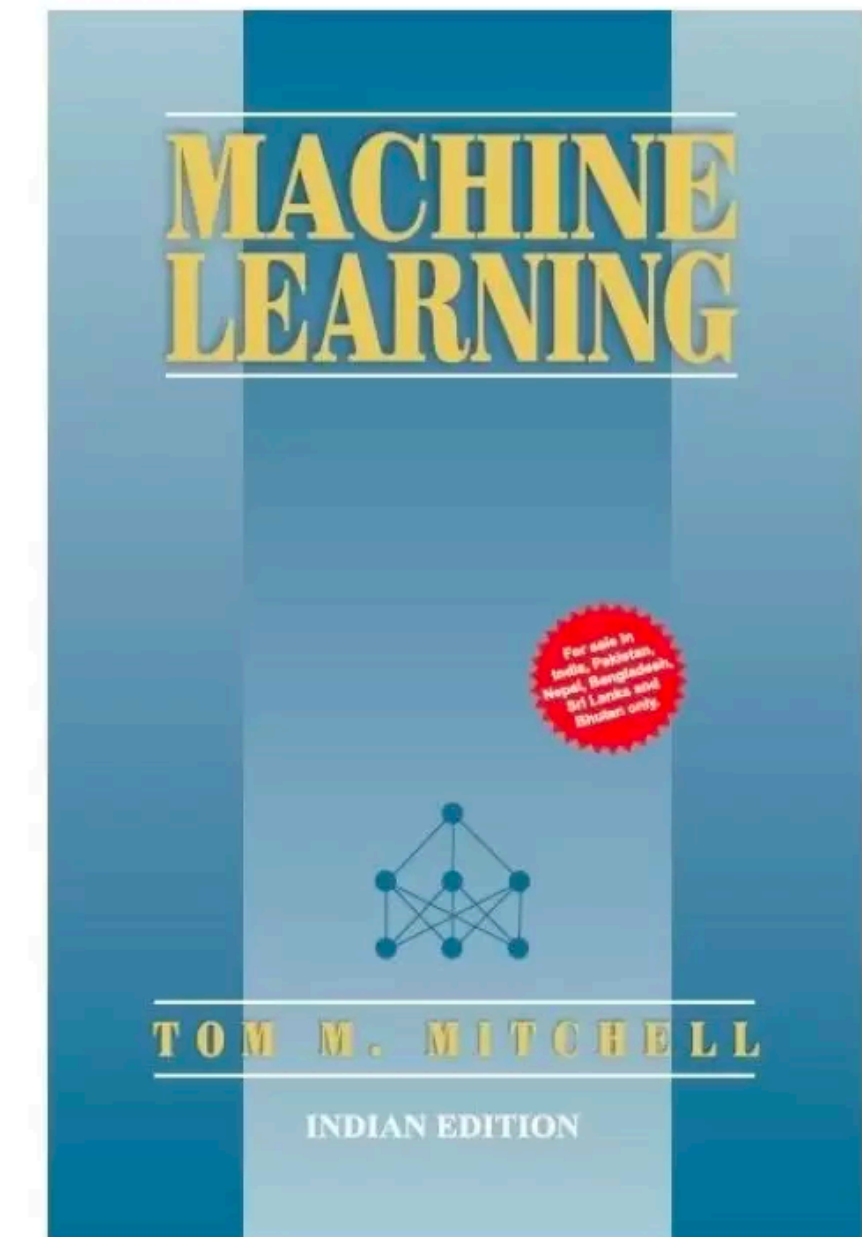
UNIVERSITY OF
TORONTO

Machine Learning:

Study of algorithms that

- improve their performance P
- at some task T
- with experience E

well-defined learning task: $\langle P, T, E \rangle$



Tom Mitchell, CMU 10-601 slides

Why should this work at all?

The main theoretical basis of supervised learning:

With a **sufficient amount of “similar” data**

+

an **expressive model class**:

Minimizing the loss function on the training data yields a highly **accurate model on unseen test data**, with high probability

1. **Data**: $S = \{(x_i, y_i)\}_{i=1, \dots, n}$

- x_i : data example with d attributes
- y_i : label of example (what you care about)

2. **Classification model**: a function $f_{(a,b,c,\dots)}$

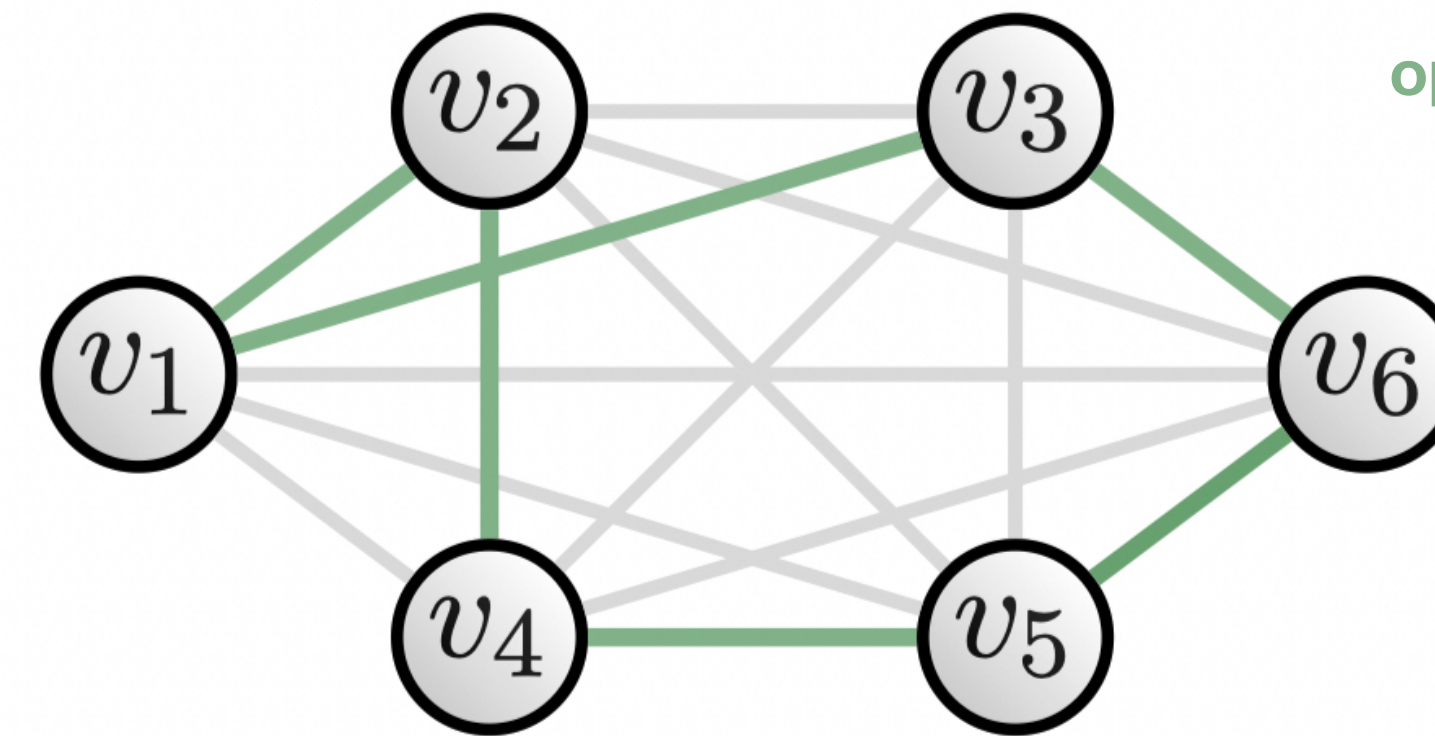
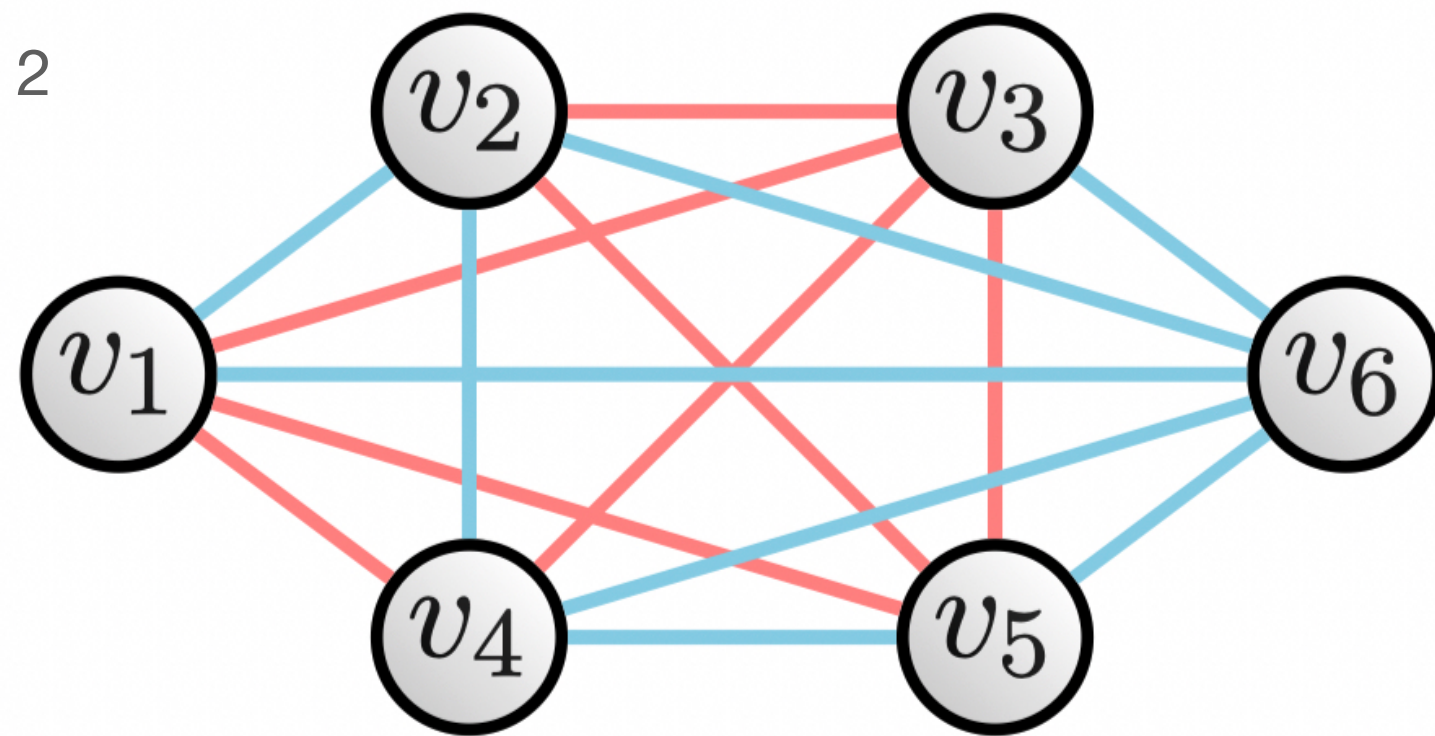
- Maps from X to Y
- $(a,b,c,\dots)$ are the **parameters**

3. **Loss function**: $L(y, f(x))$

- Penalizes the model's mistakes

Graph Optimization

— cost = 1
— cost = 2

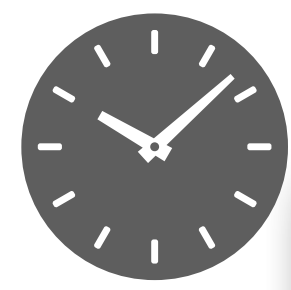


optimal solution, cost = 7

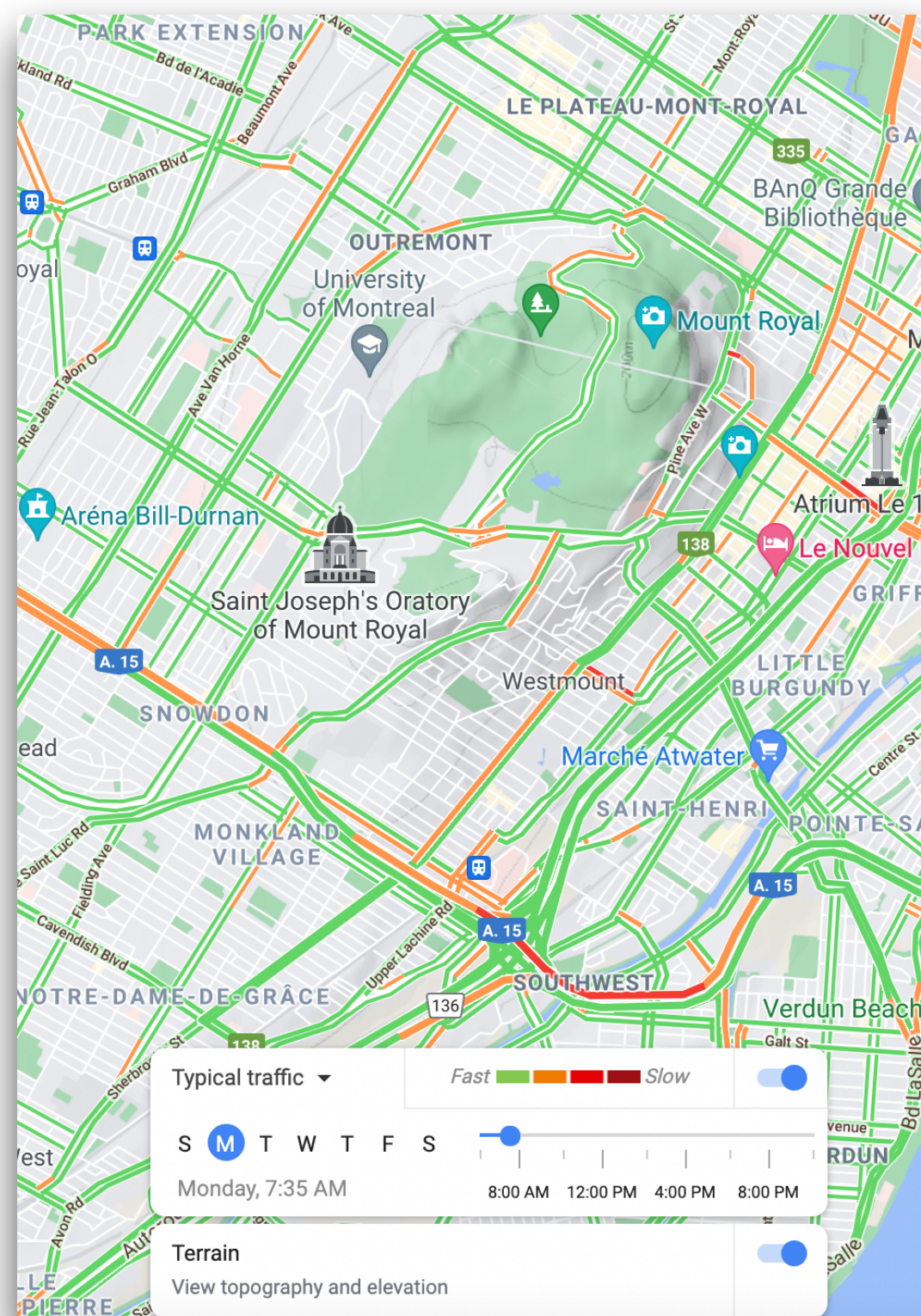
Travelling Salesperson Problem (TSP)

TSPs in Montréal

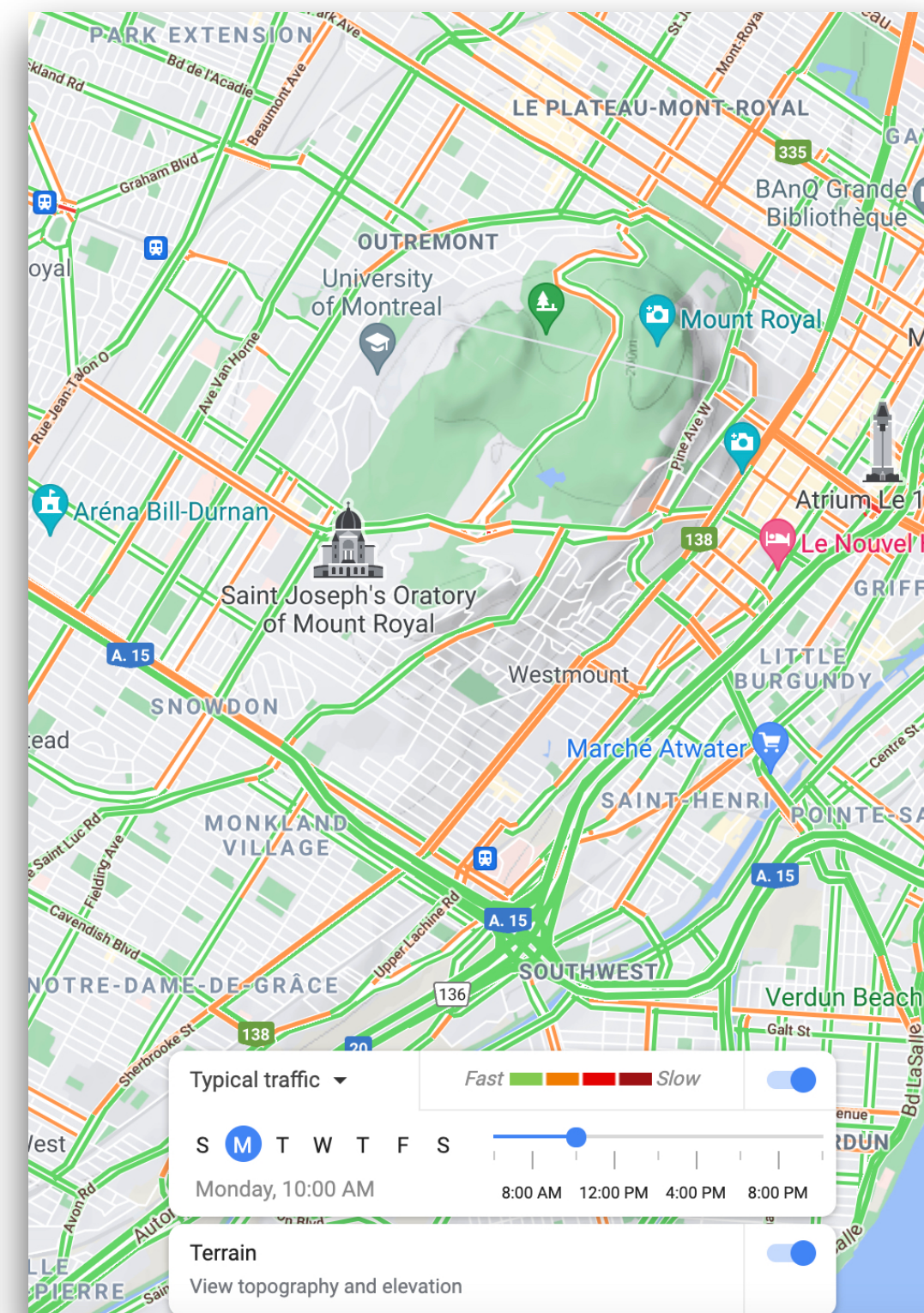
Every few hours, optimally **route** the salesperson through a set of locations



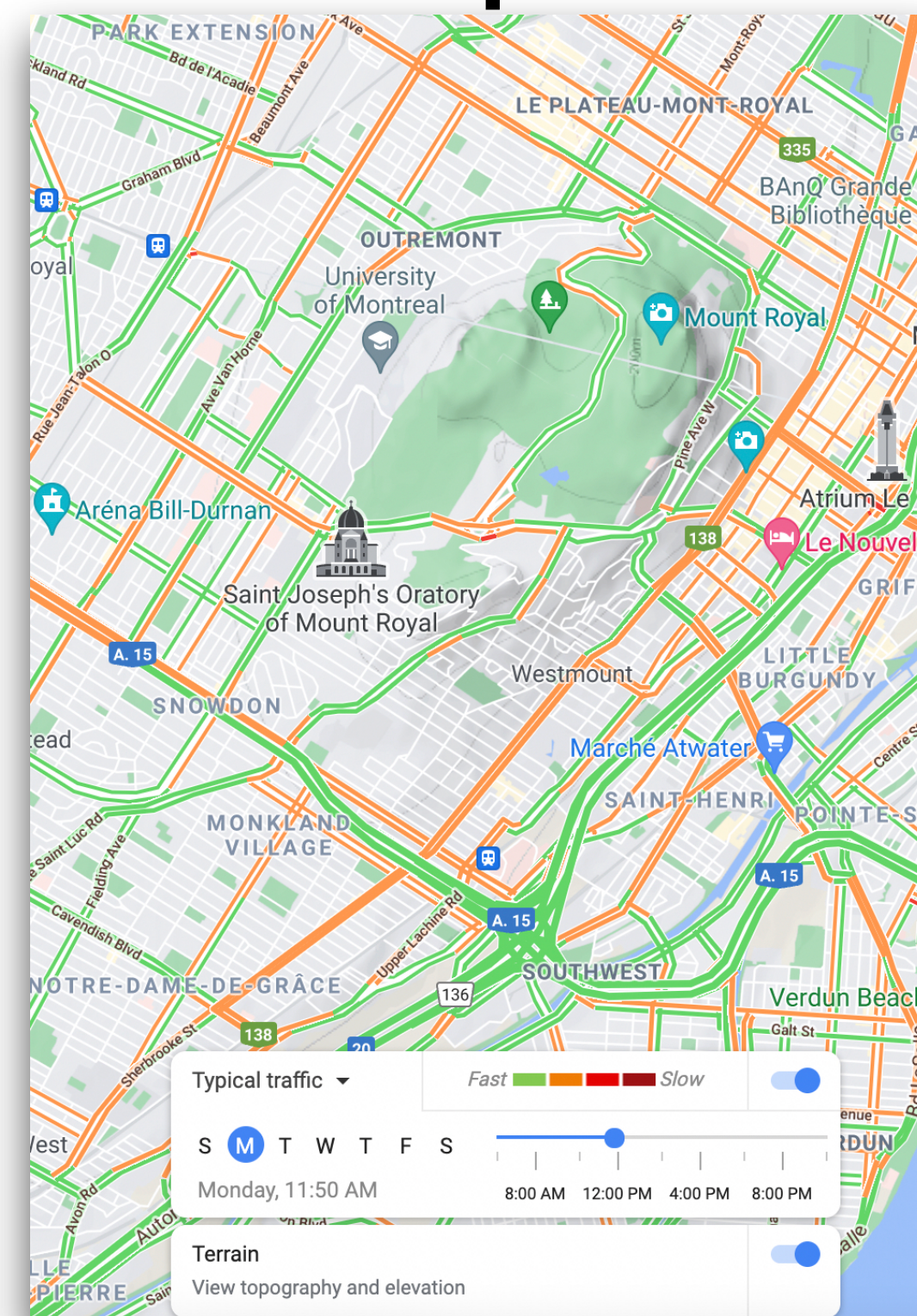
8 am



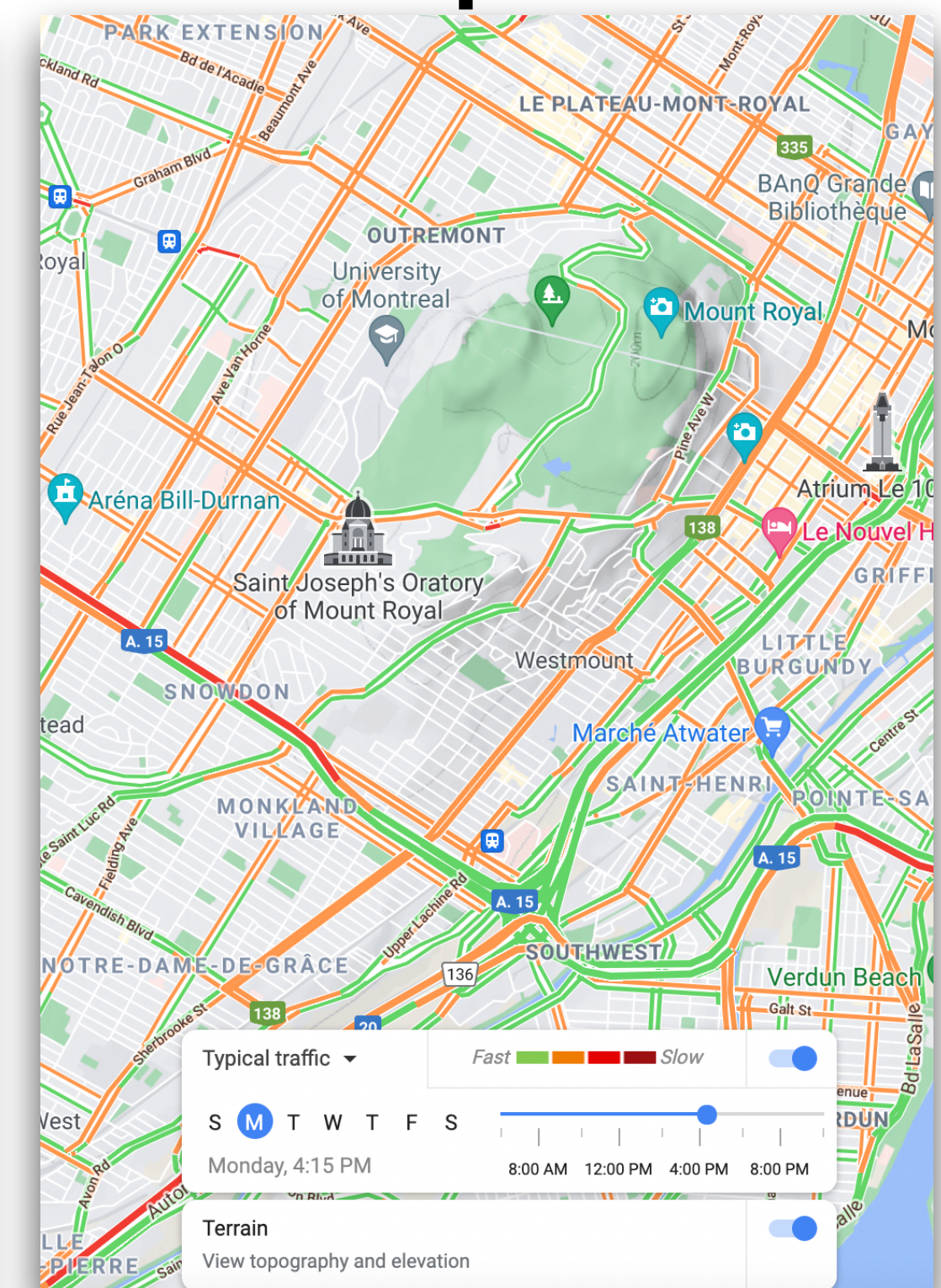
10 am



12 pm



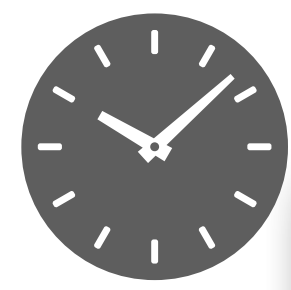
4 pm



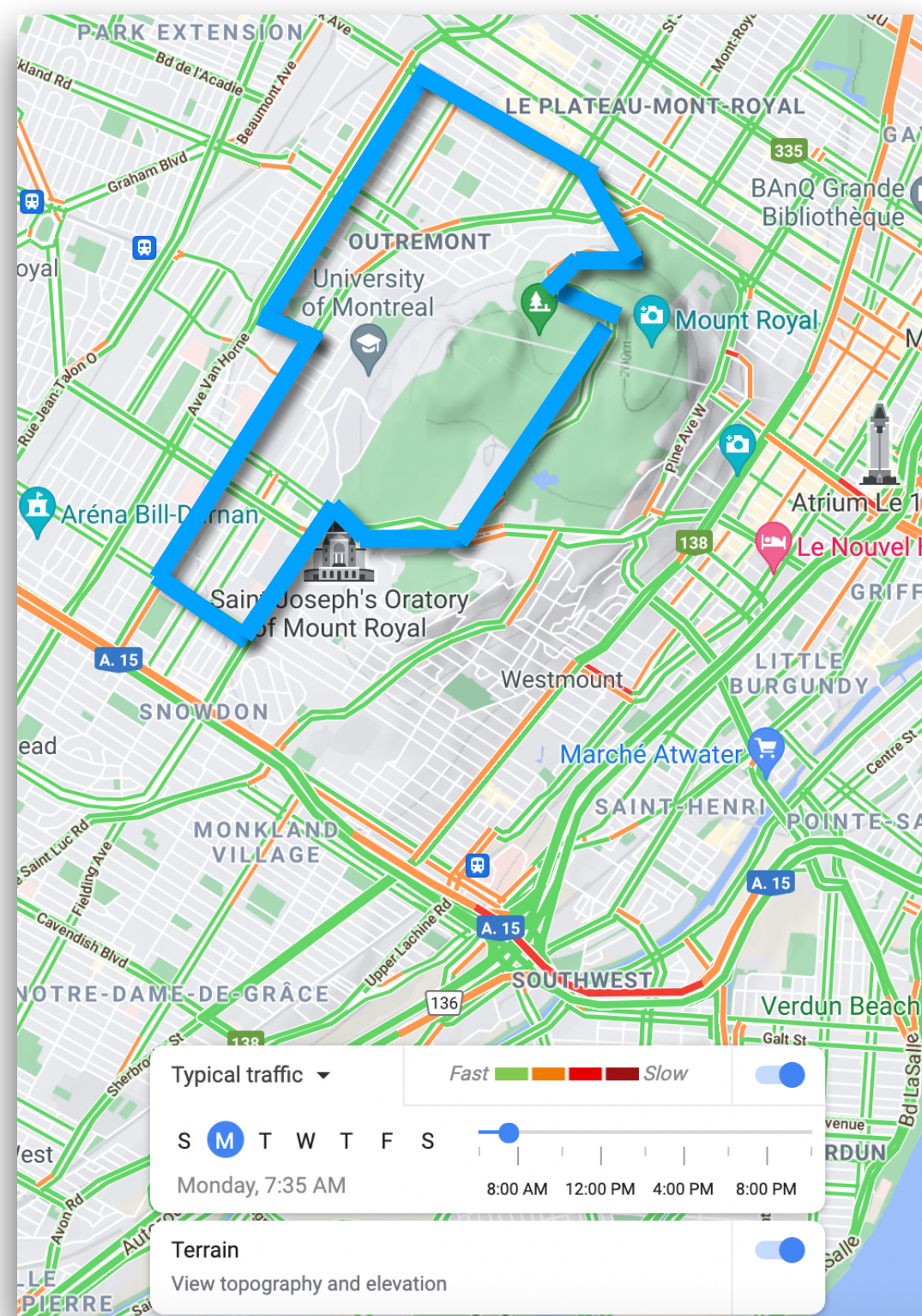
Google Maps, Montréal, Québec, Canada

TSPs in Montréal

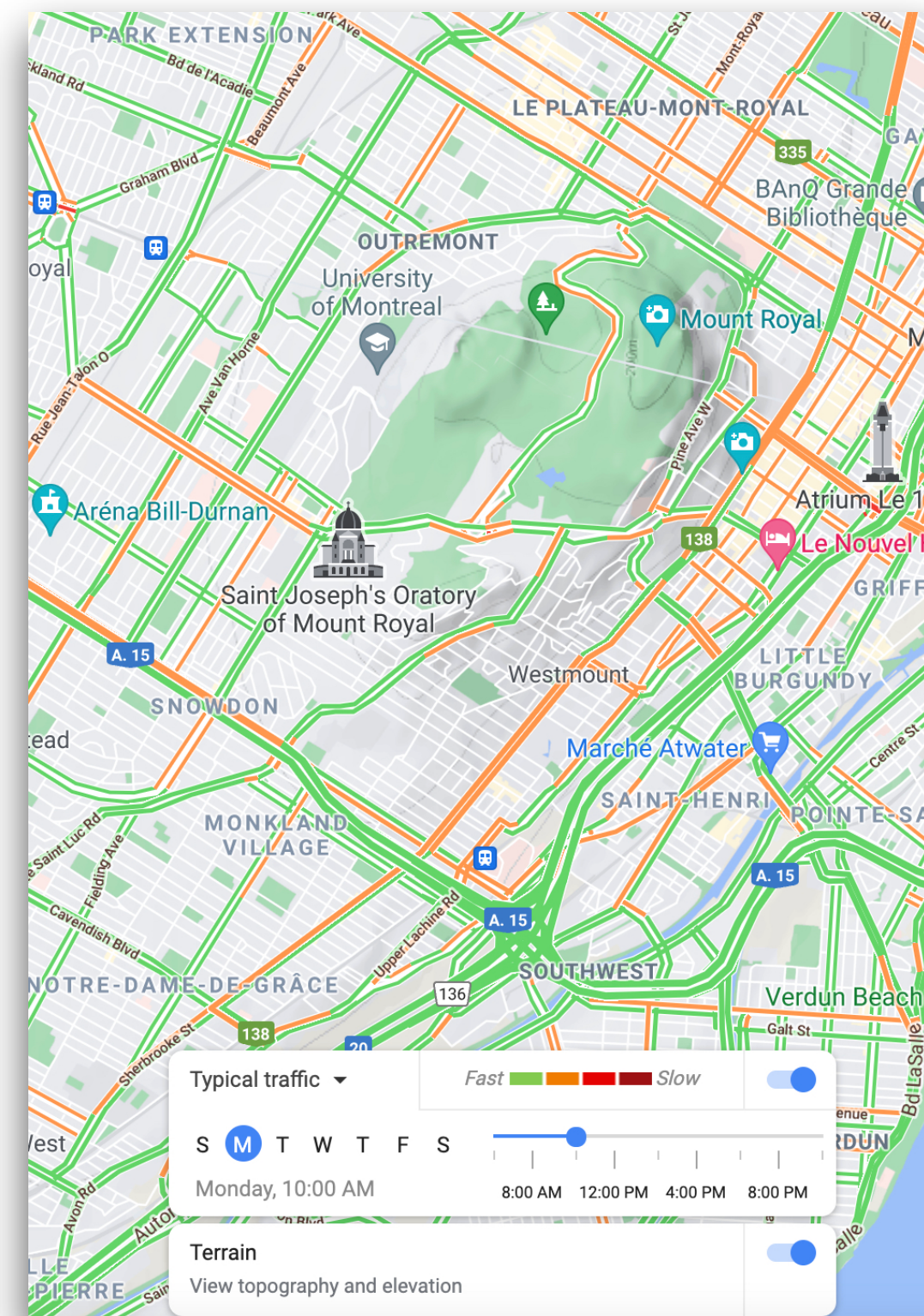
Every few hours, optimally **route** the salesperson through a set of locations



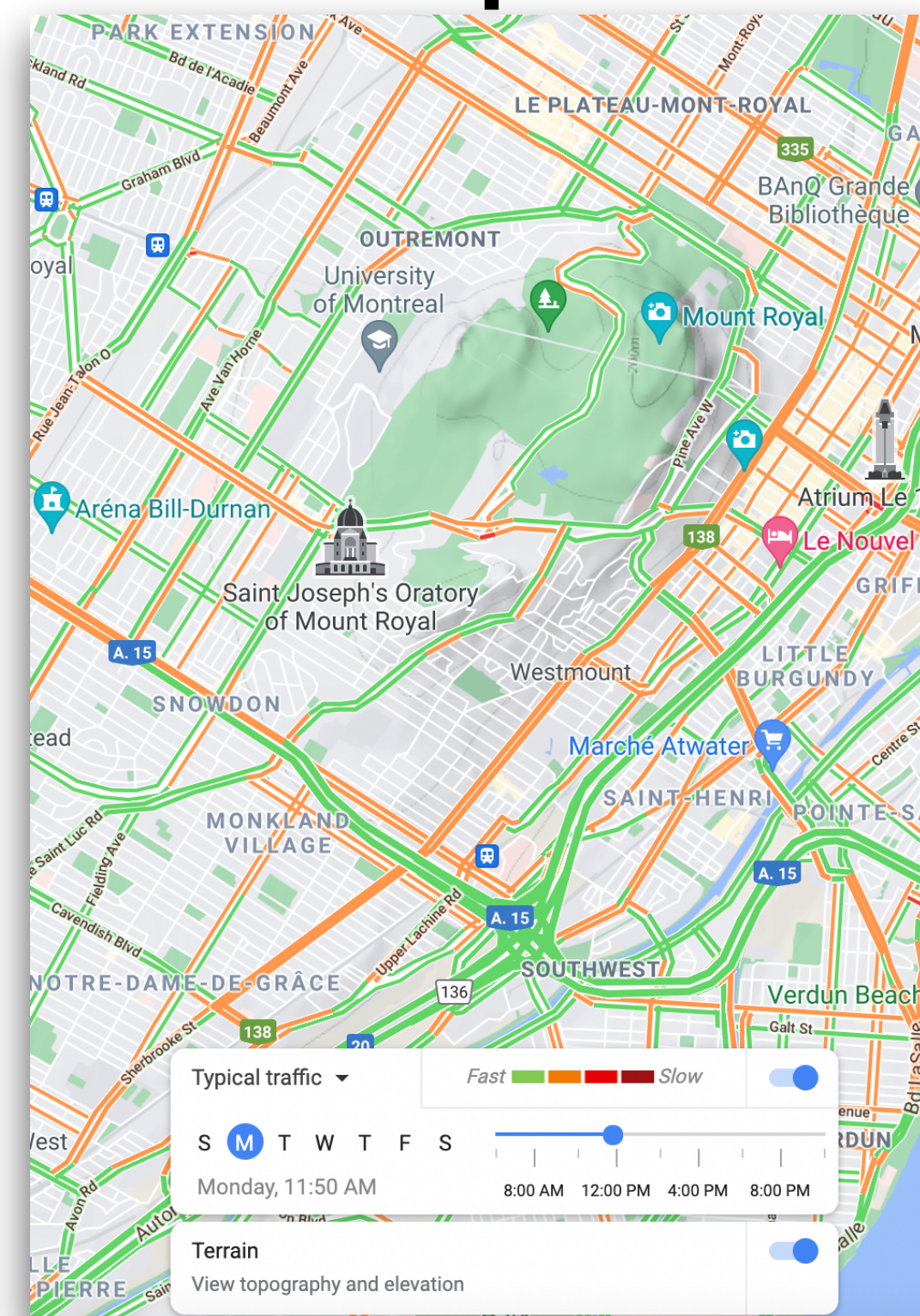
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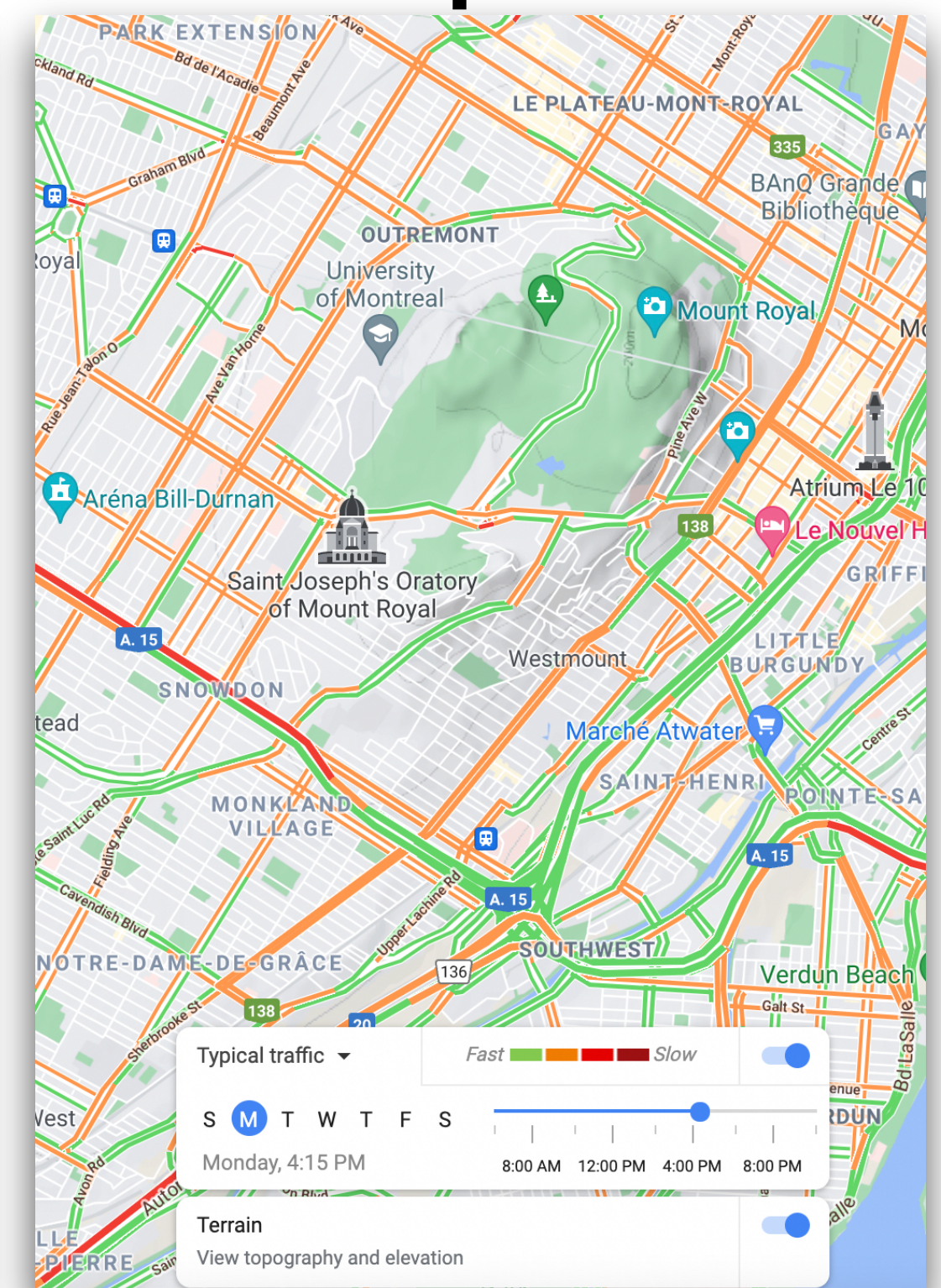
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12 pm



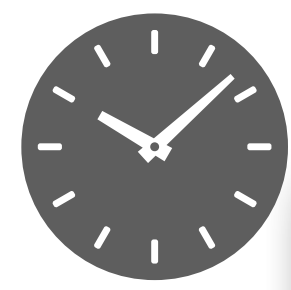
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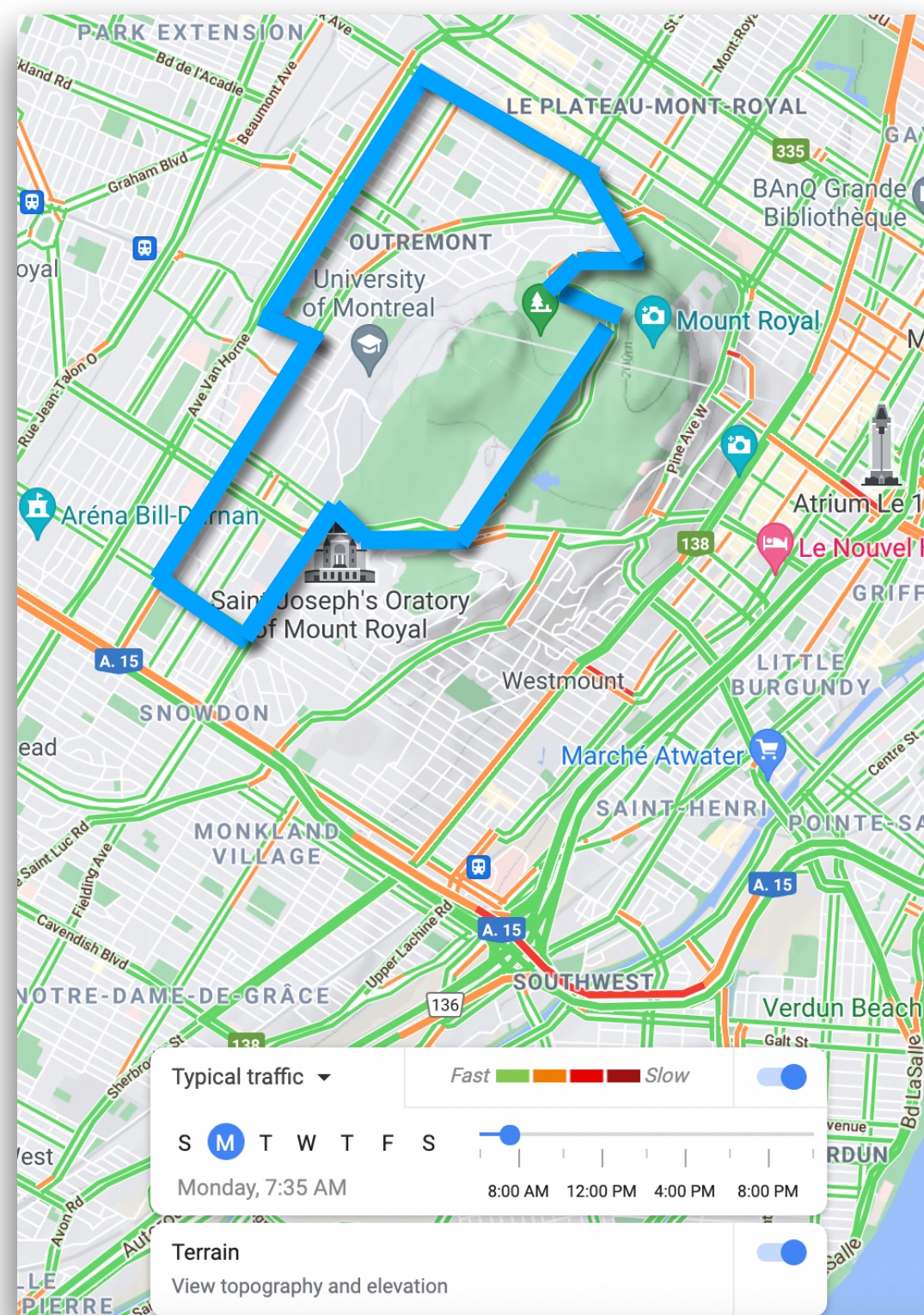
Google Maps, Montréal, Québec, Canada

TSPs in Montréal

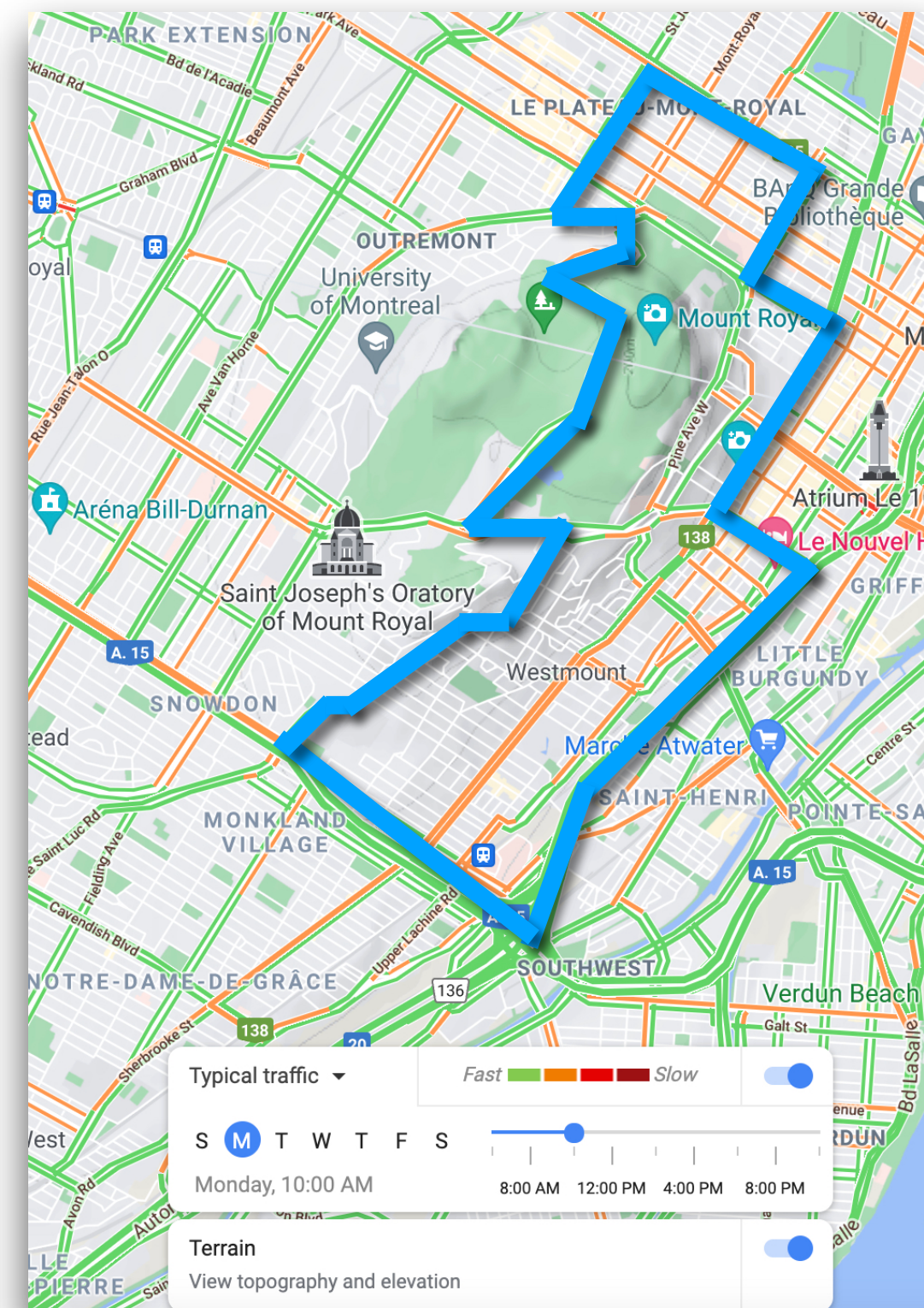
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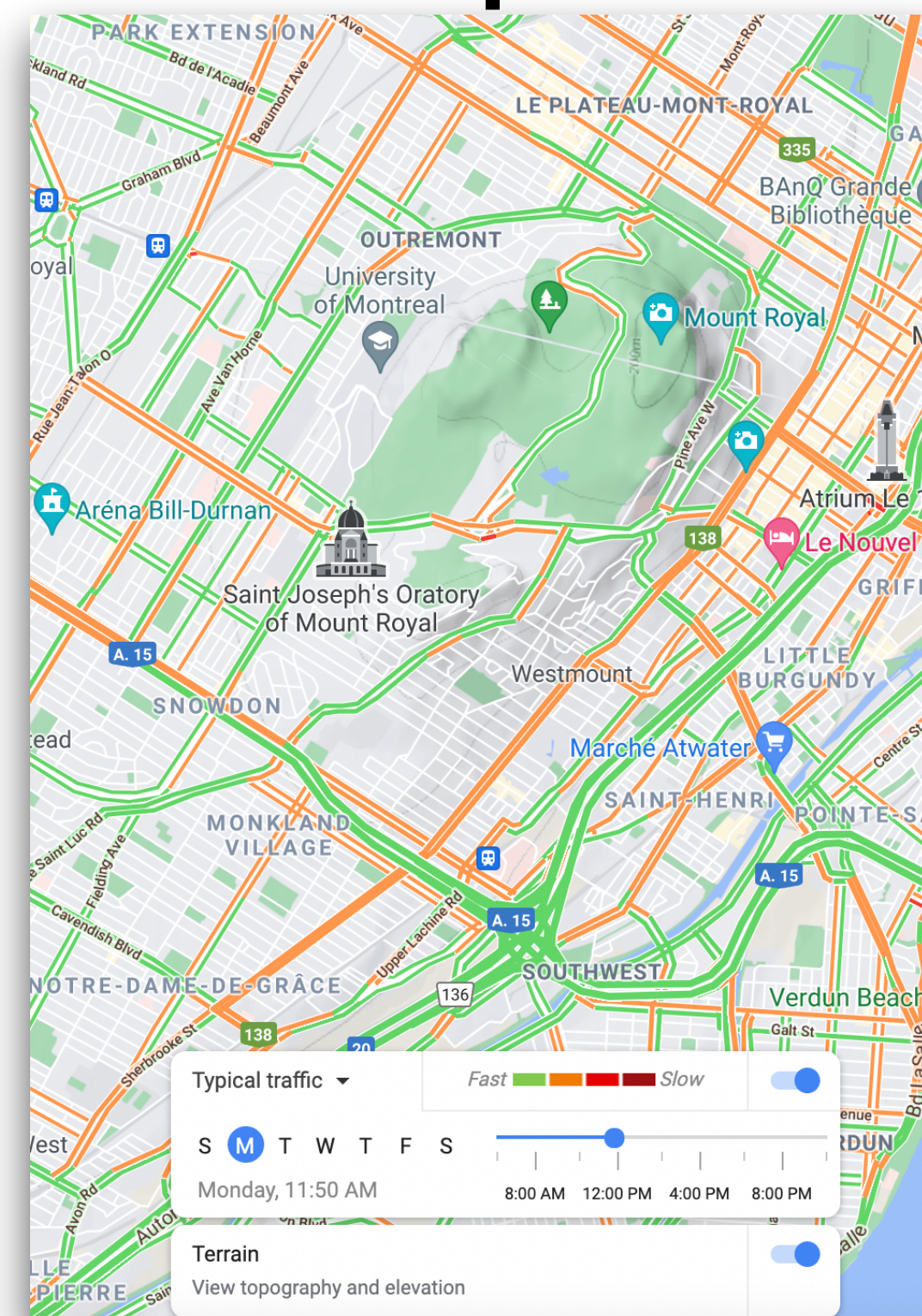
8 am



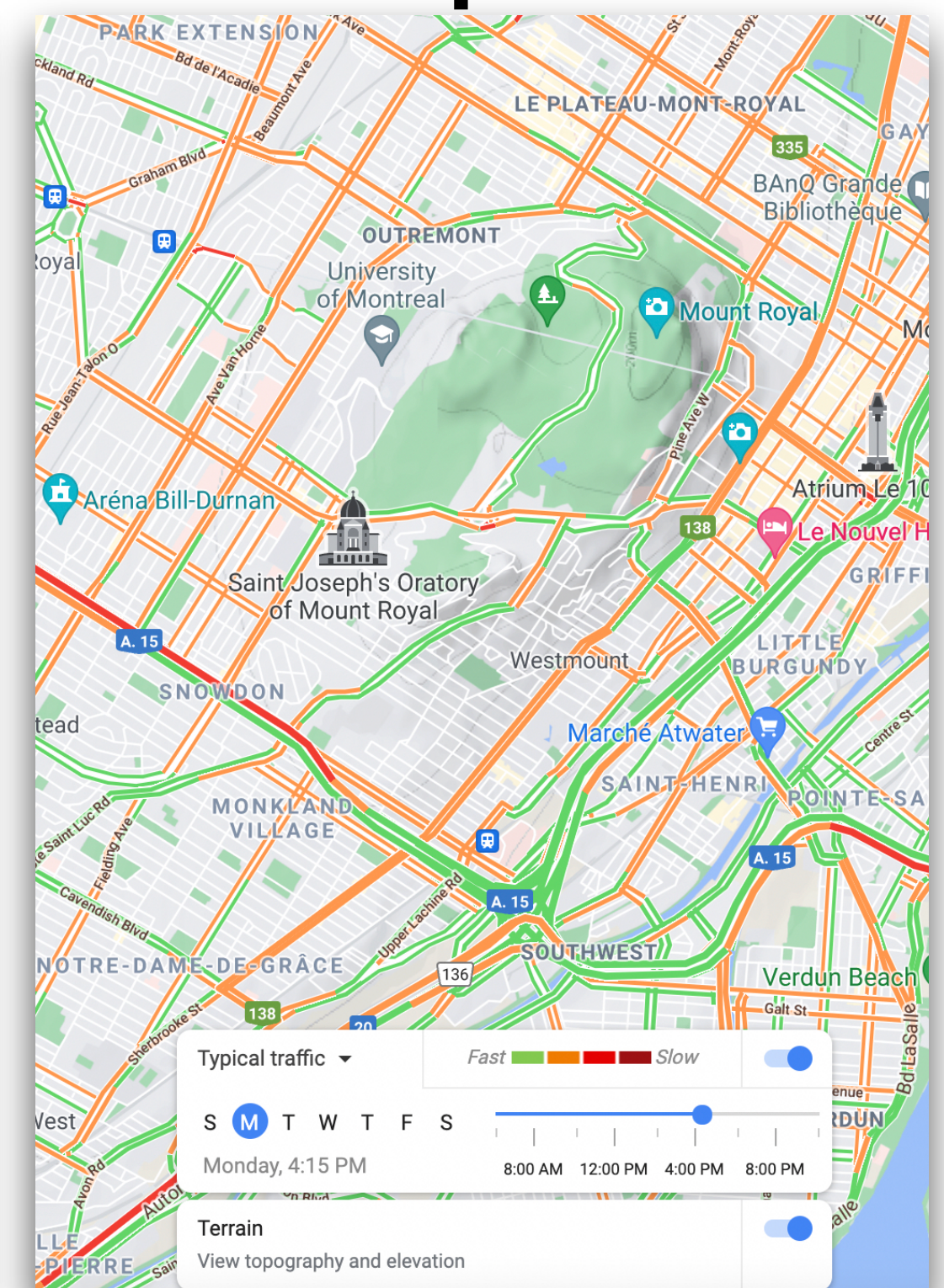
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12 pm



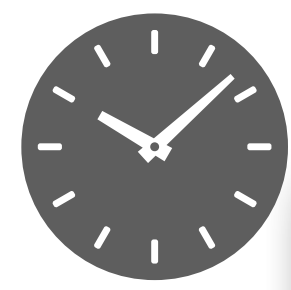
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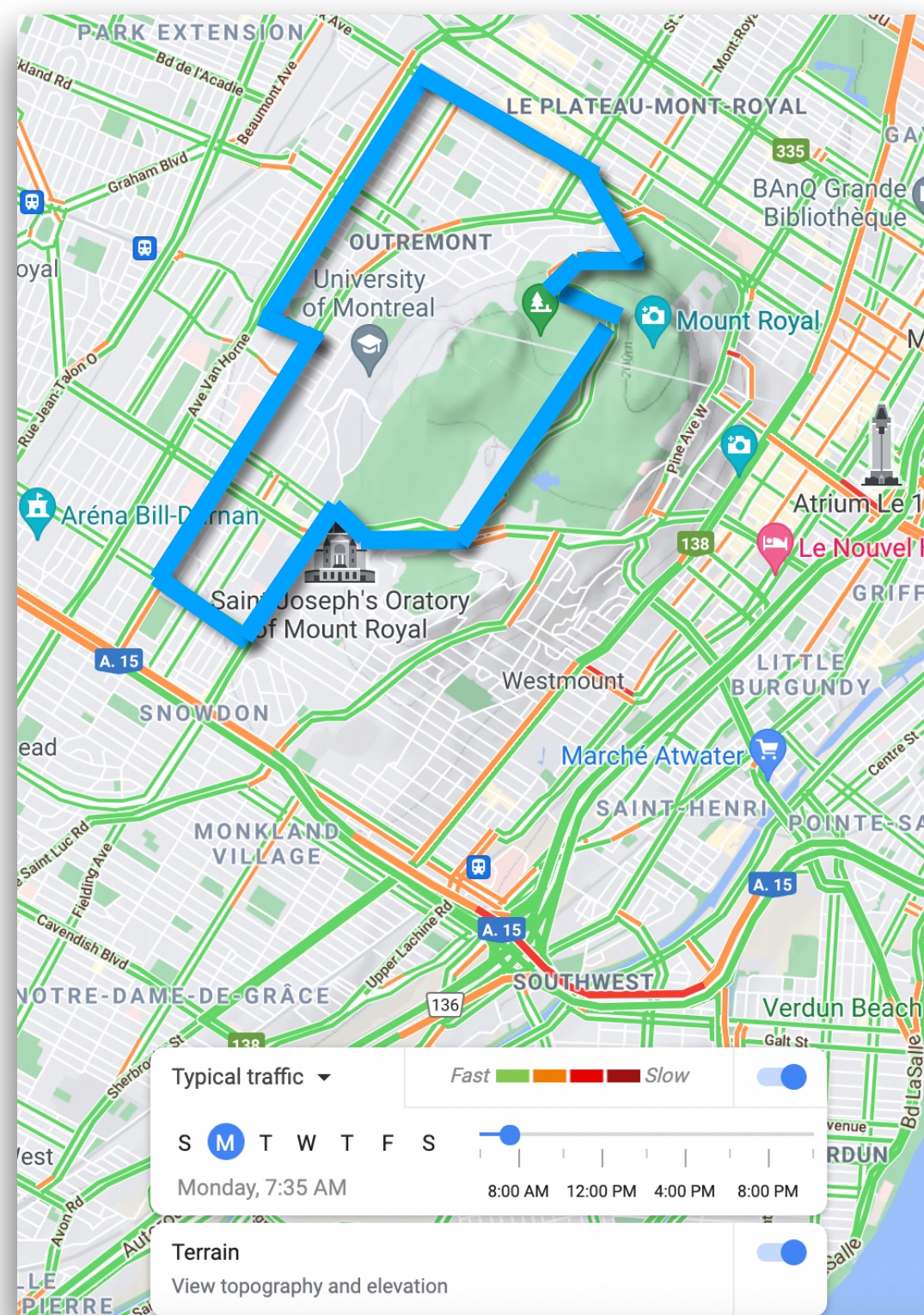
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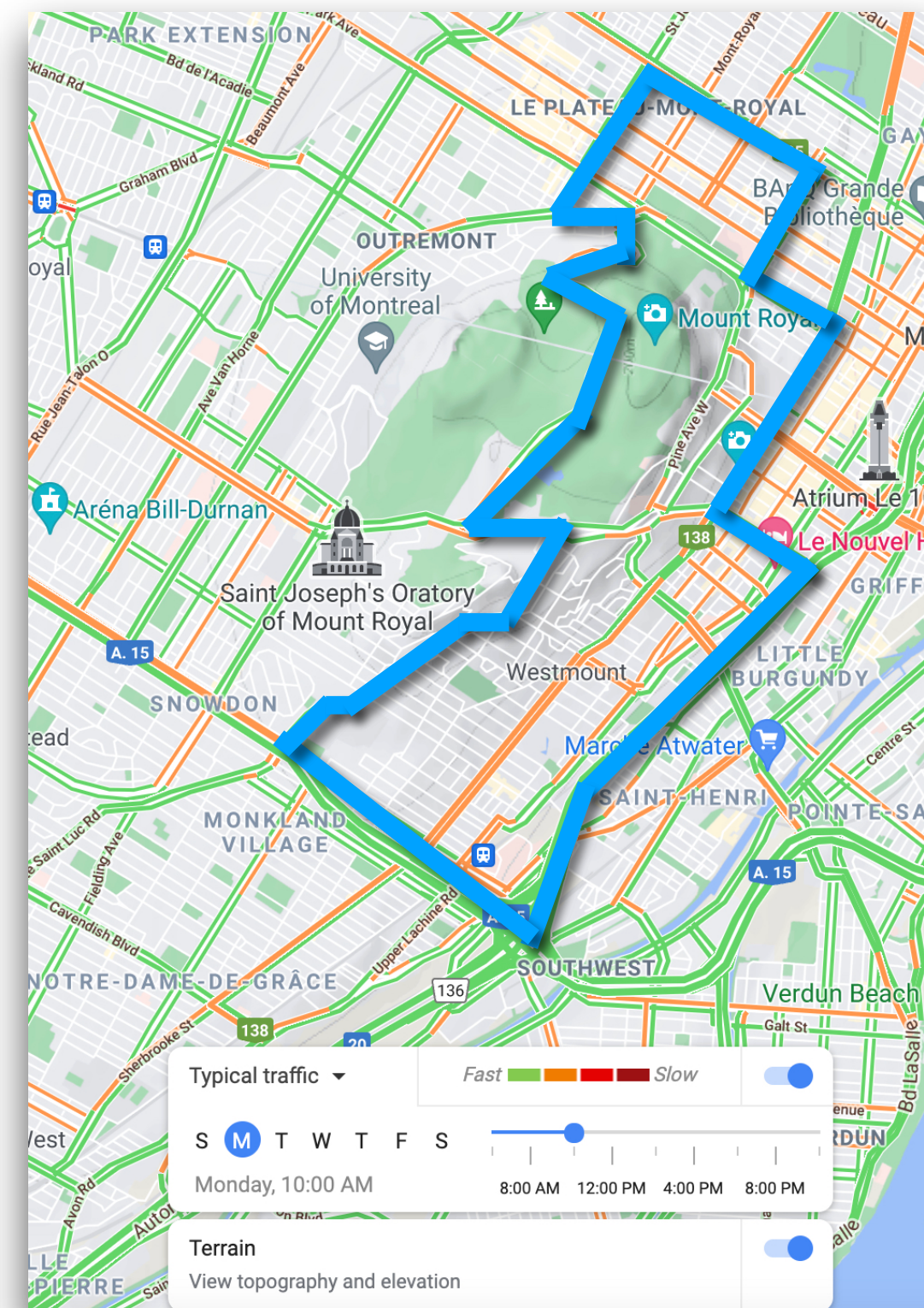
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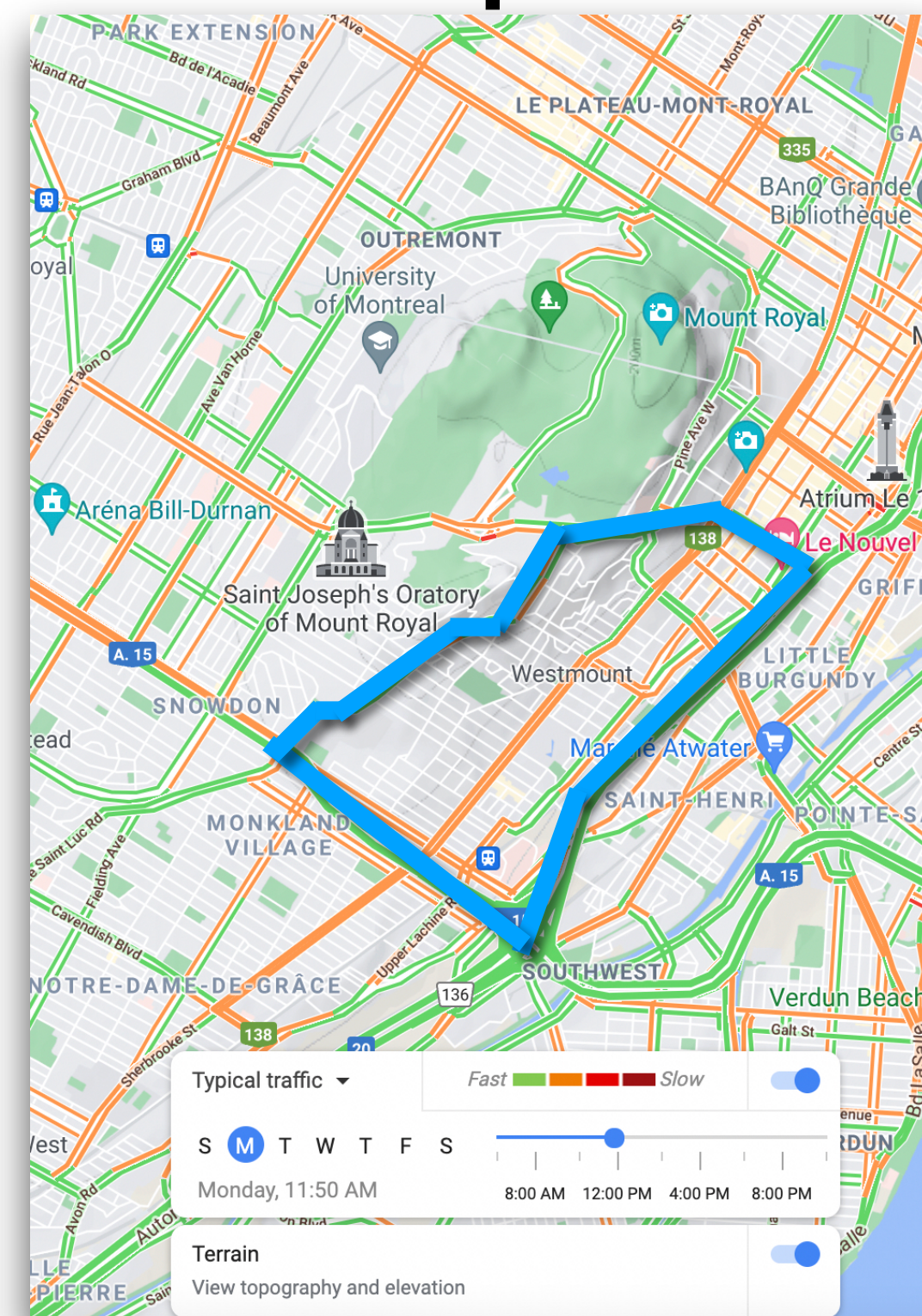
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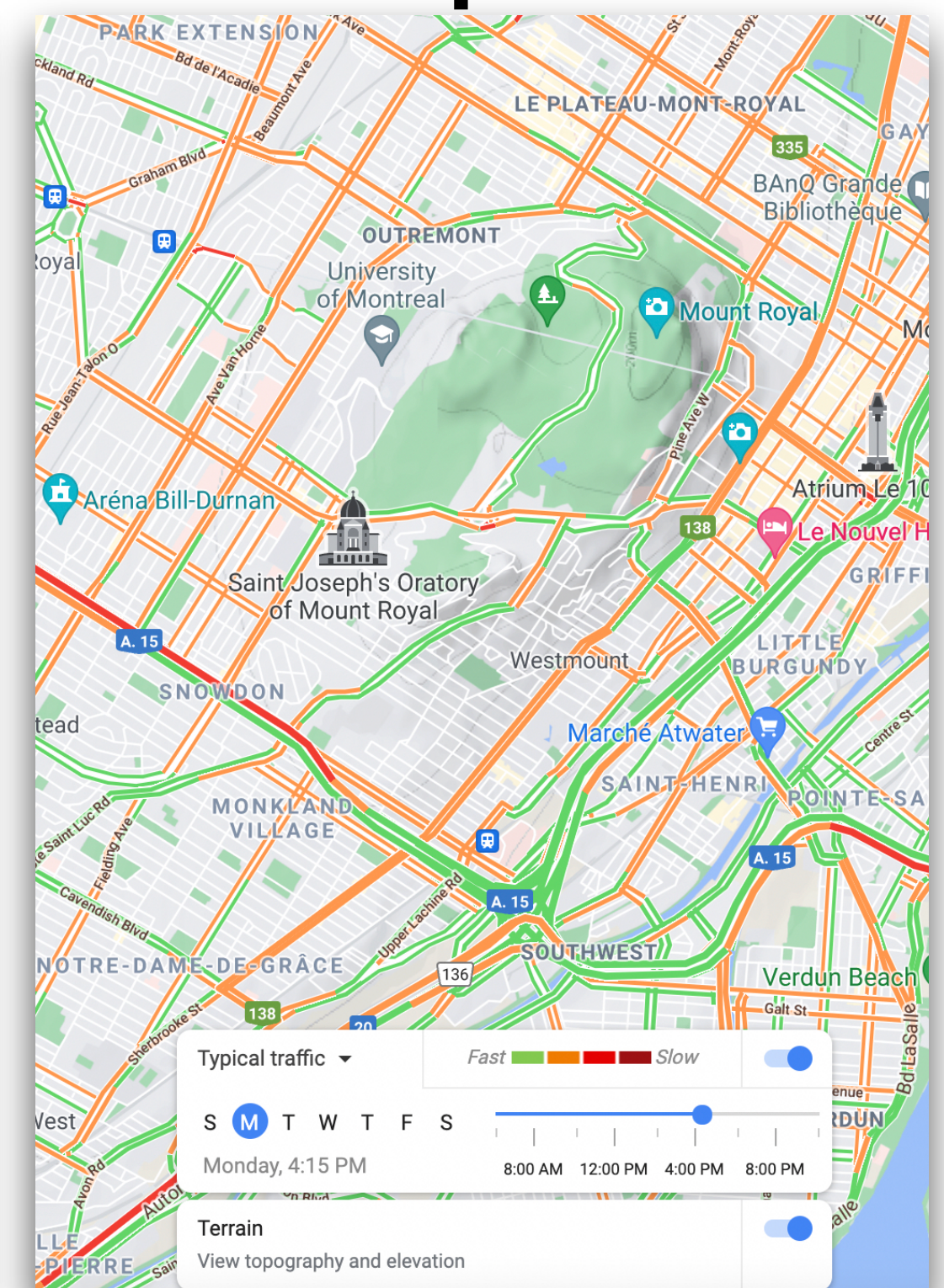
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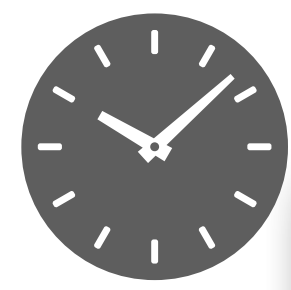
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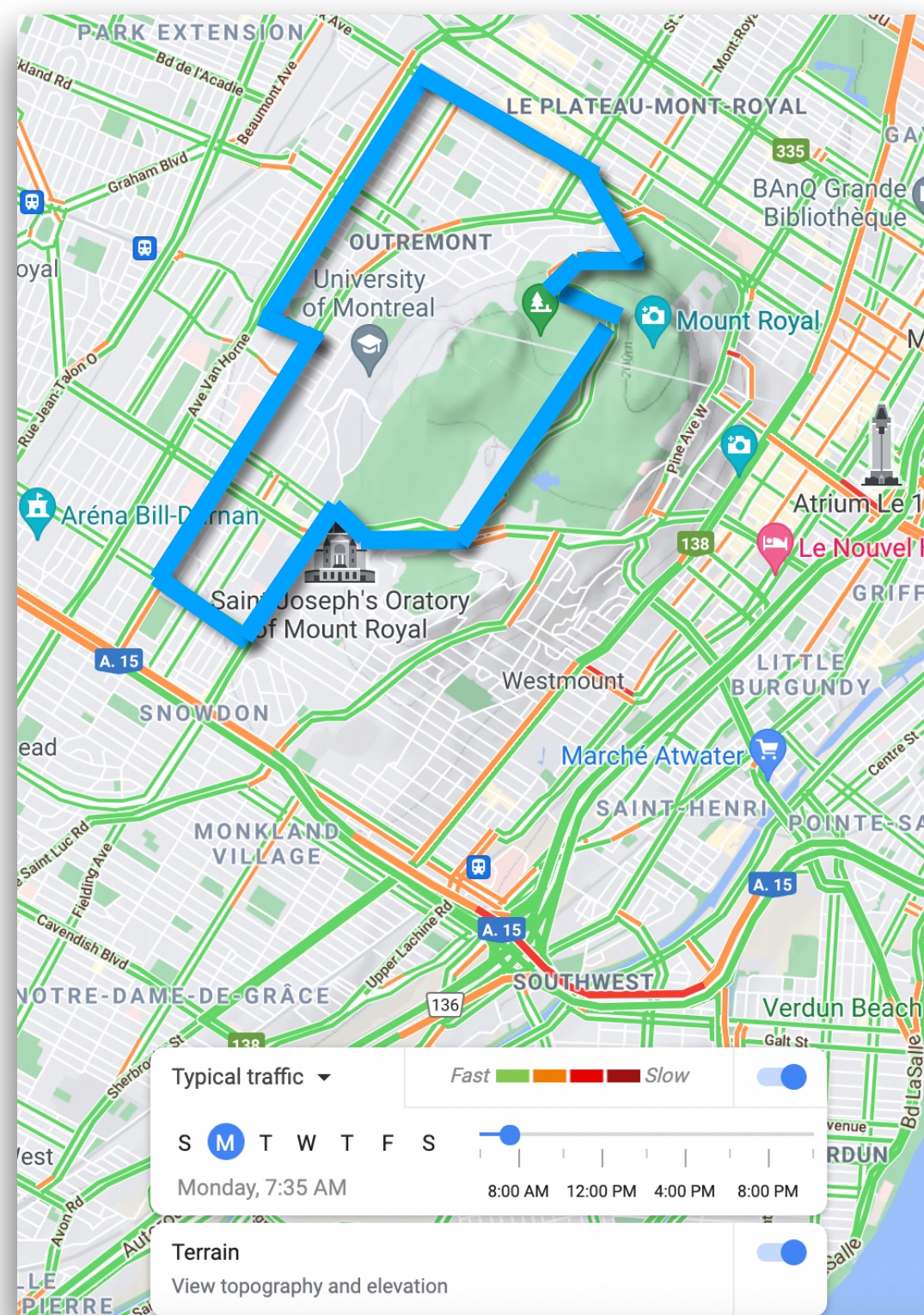
Google Maps, Montréal, Québec, Canada

TSPs in Montréal

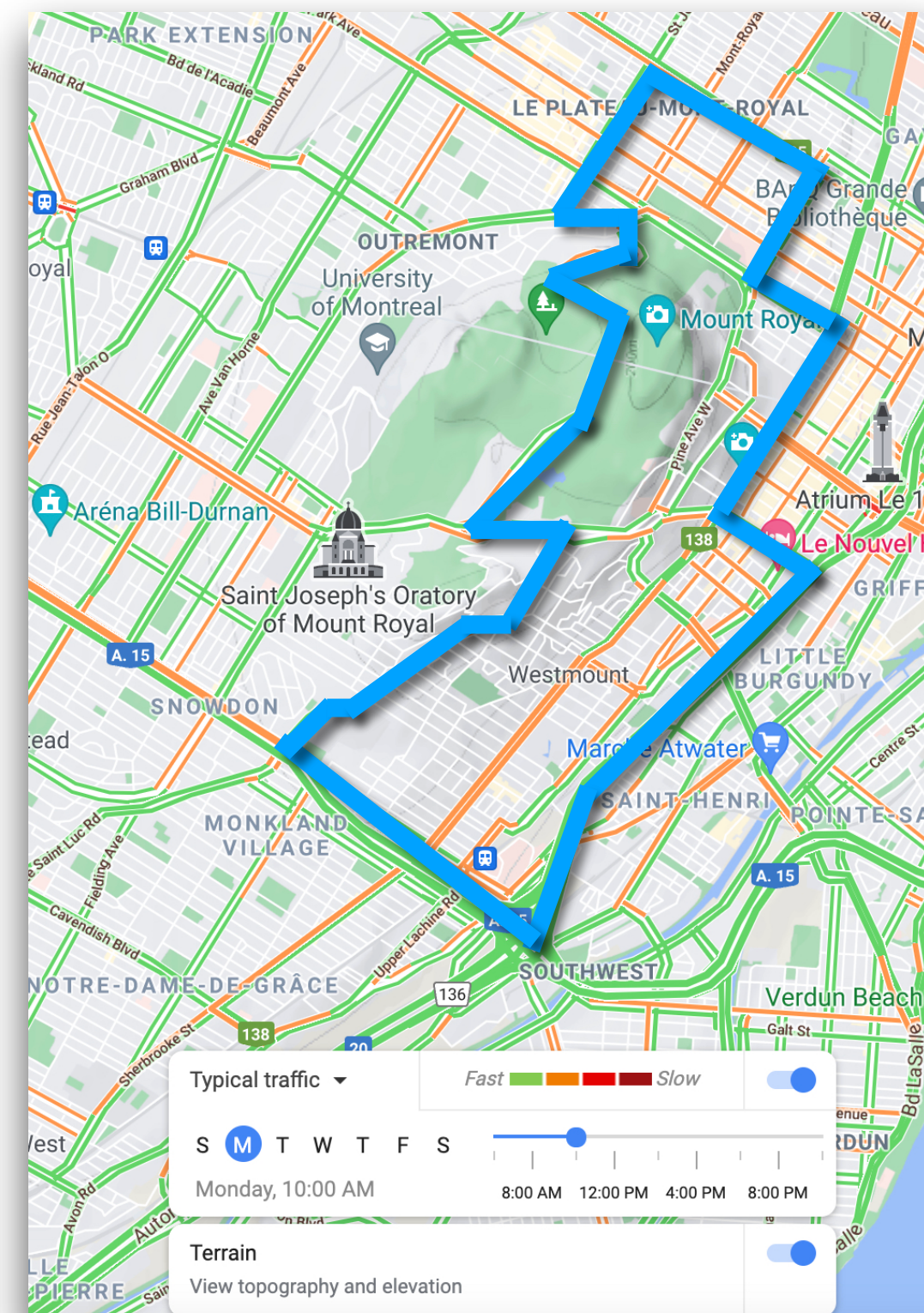
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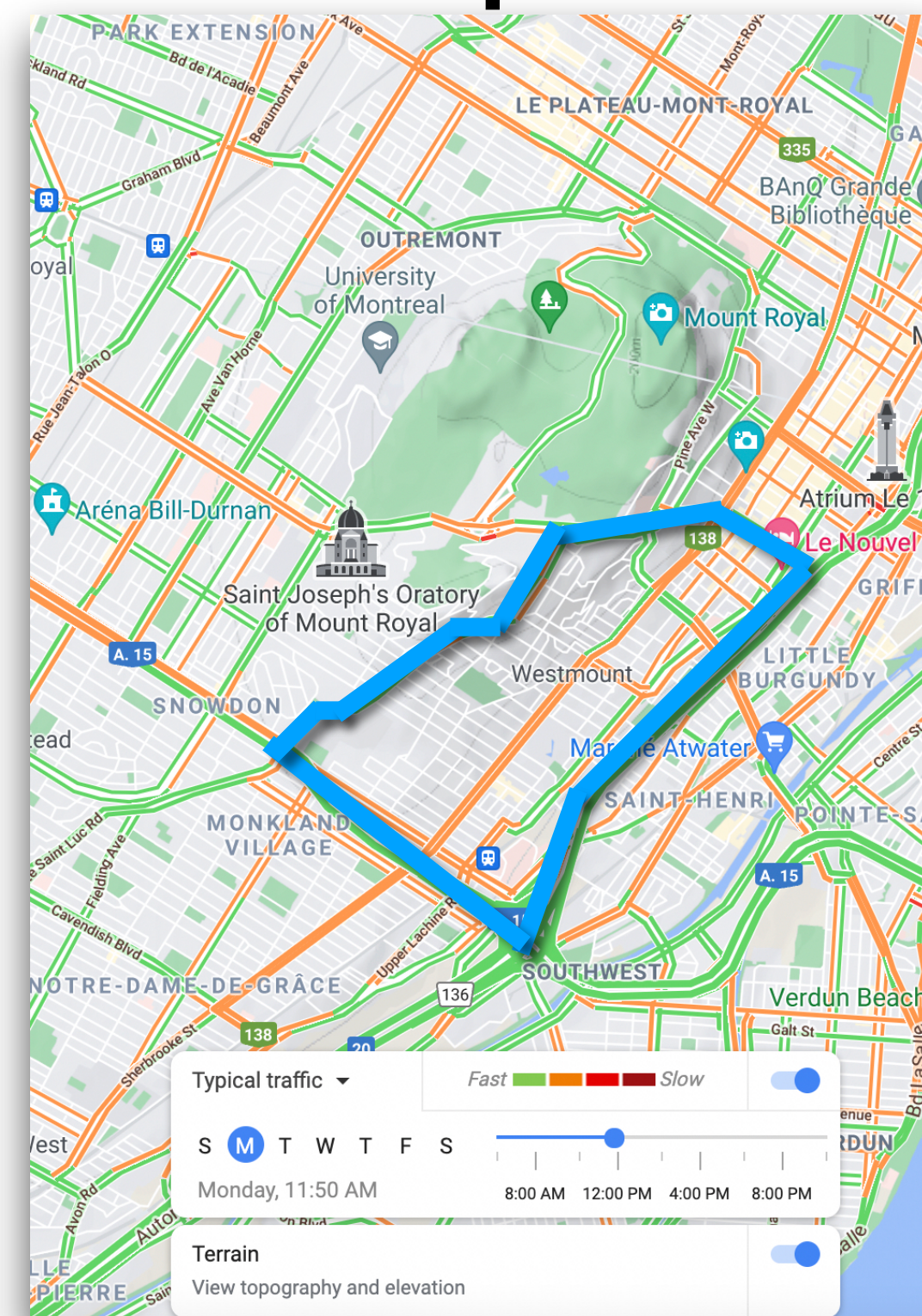
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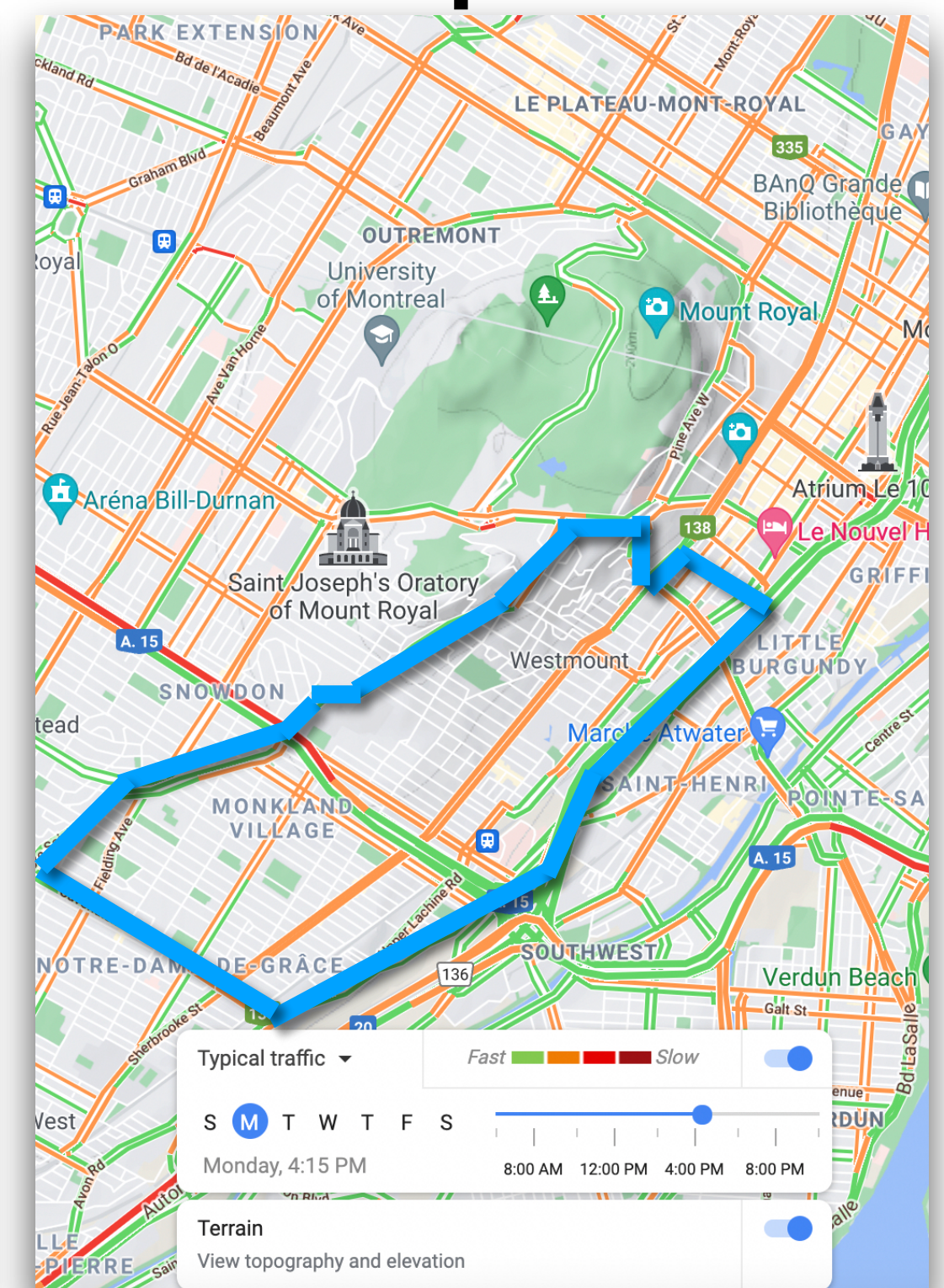
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12 pm



4 pm



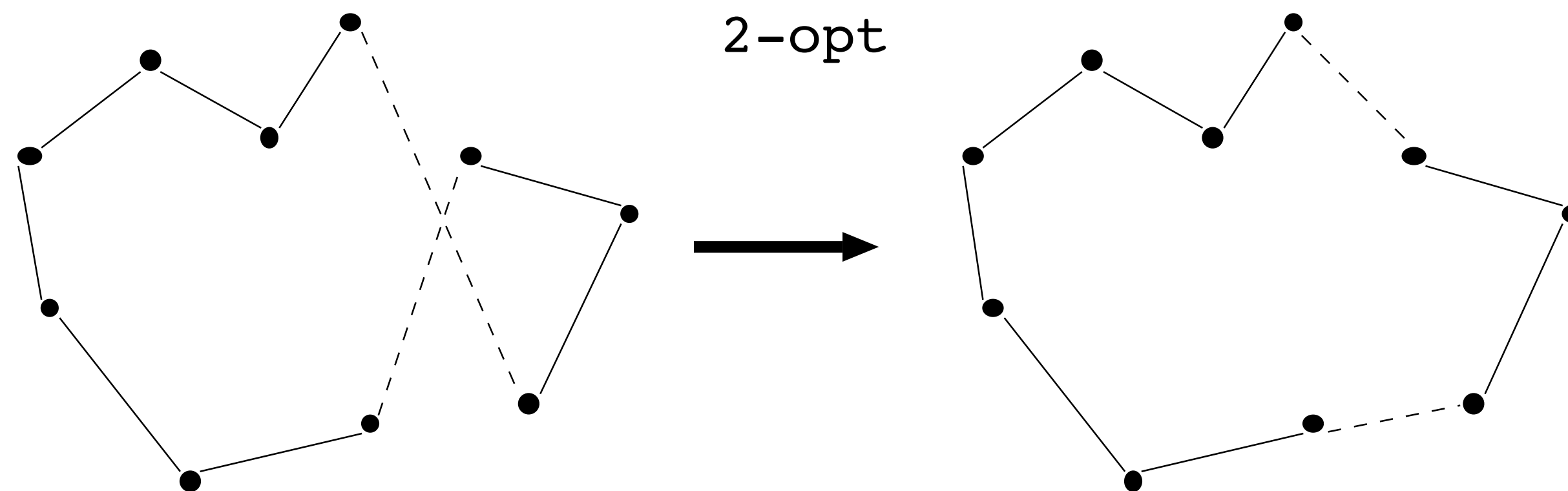
Google Maps, Montréal, Québec, Canada

Nearest Neighbour heuristic for the TSP:

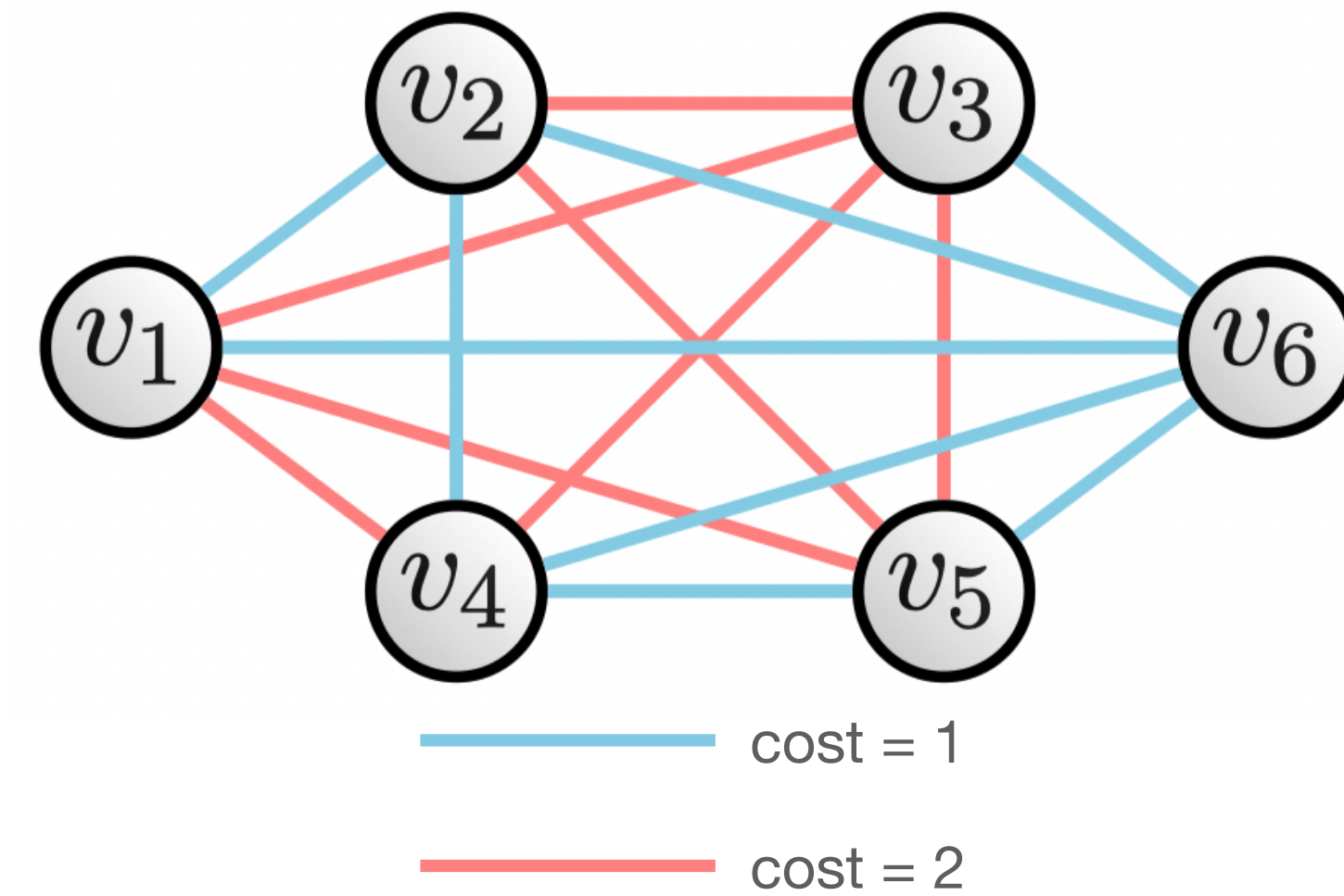
- always choose at the current city the closest unvisited city
 - choose an arbitrary initial city $\pi(1)$
 - at the i th step choose city $\pi(i + 1)$ to be the city j that minimises $\{d(\pi(i), j)\}; j \neq \pi(k), 1 \leq k \leq i$

Iterative Improvement for the TSP

- initial solution is a complete tour
- k -opt neighbourhood: solutions which differ by at most k edges



- neighbourhood size $\mathcal{O}(n^k)$

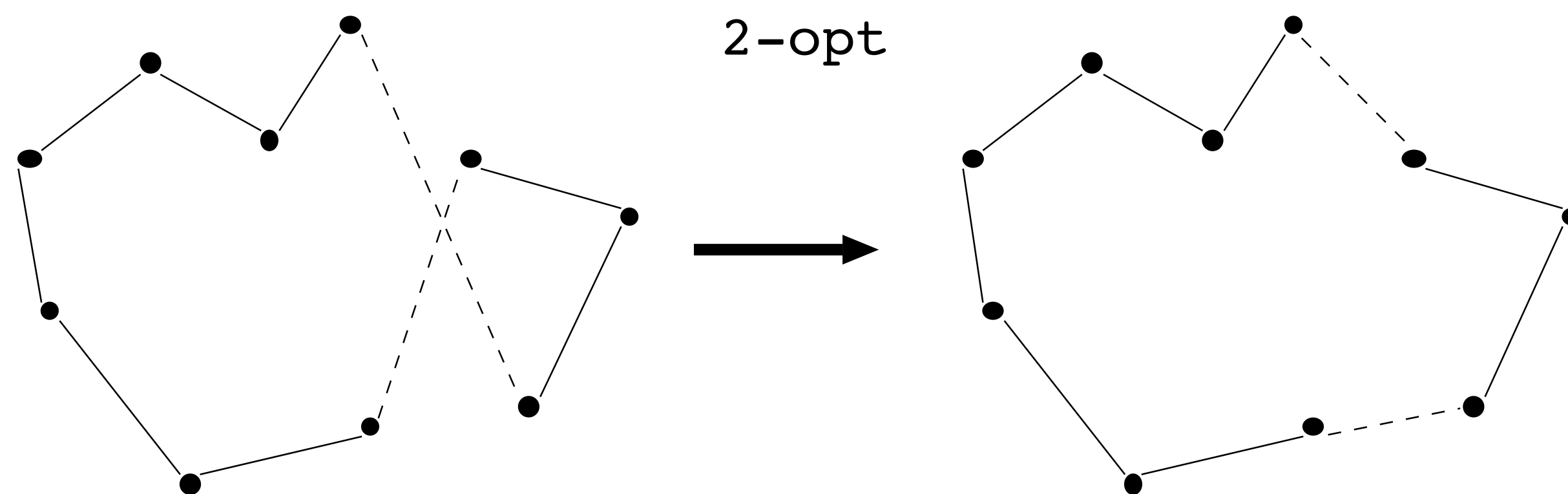


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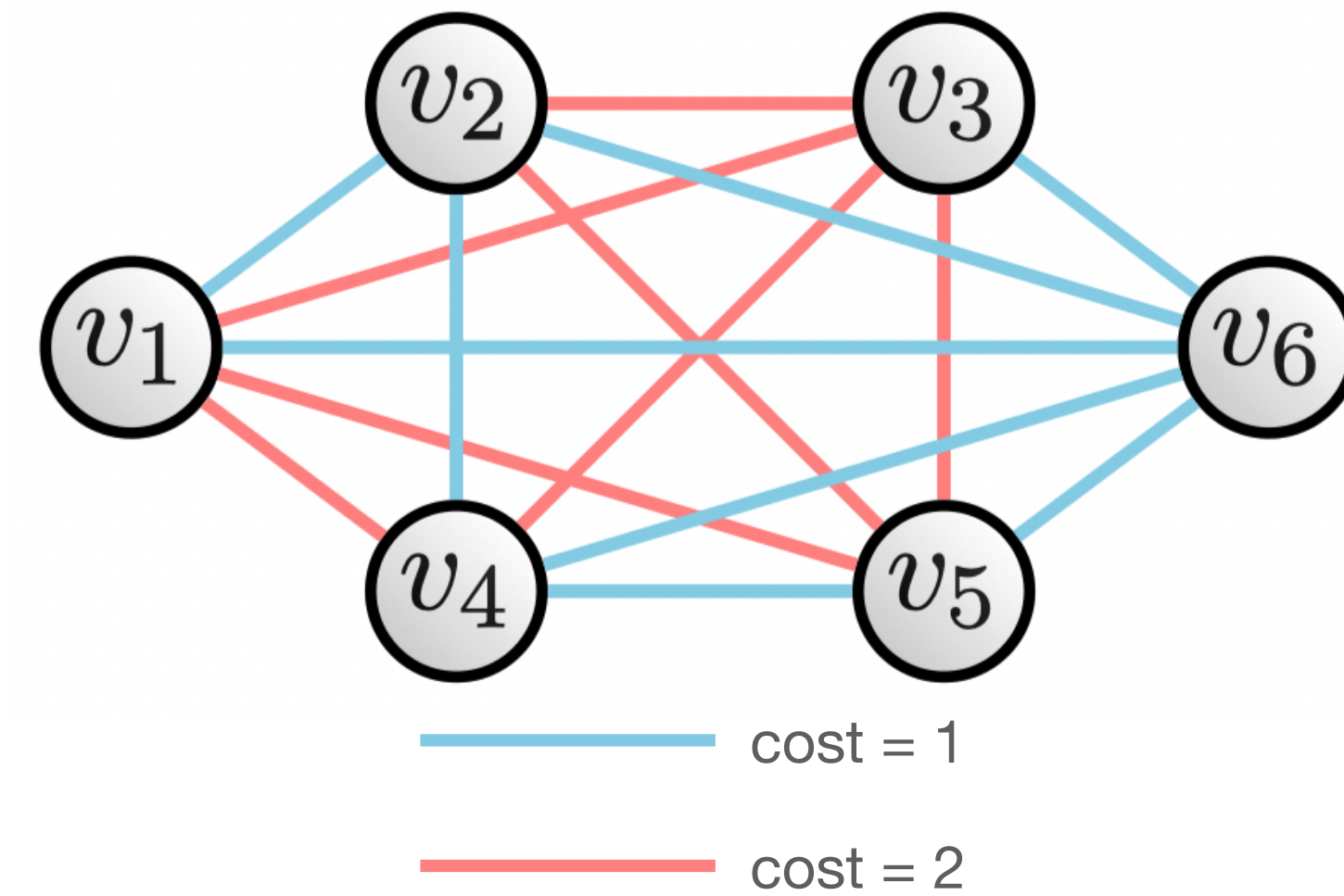
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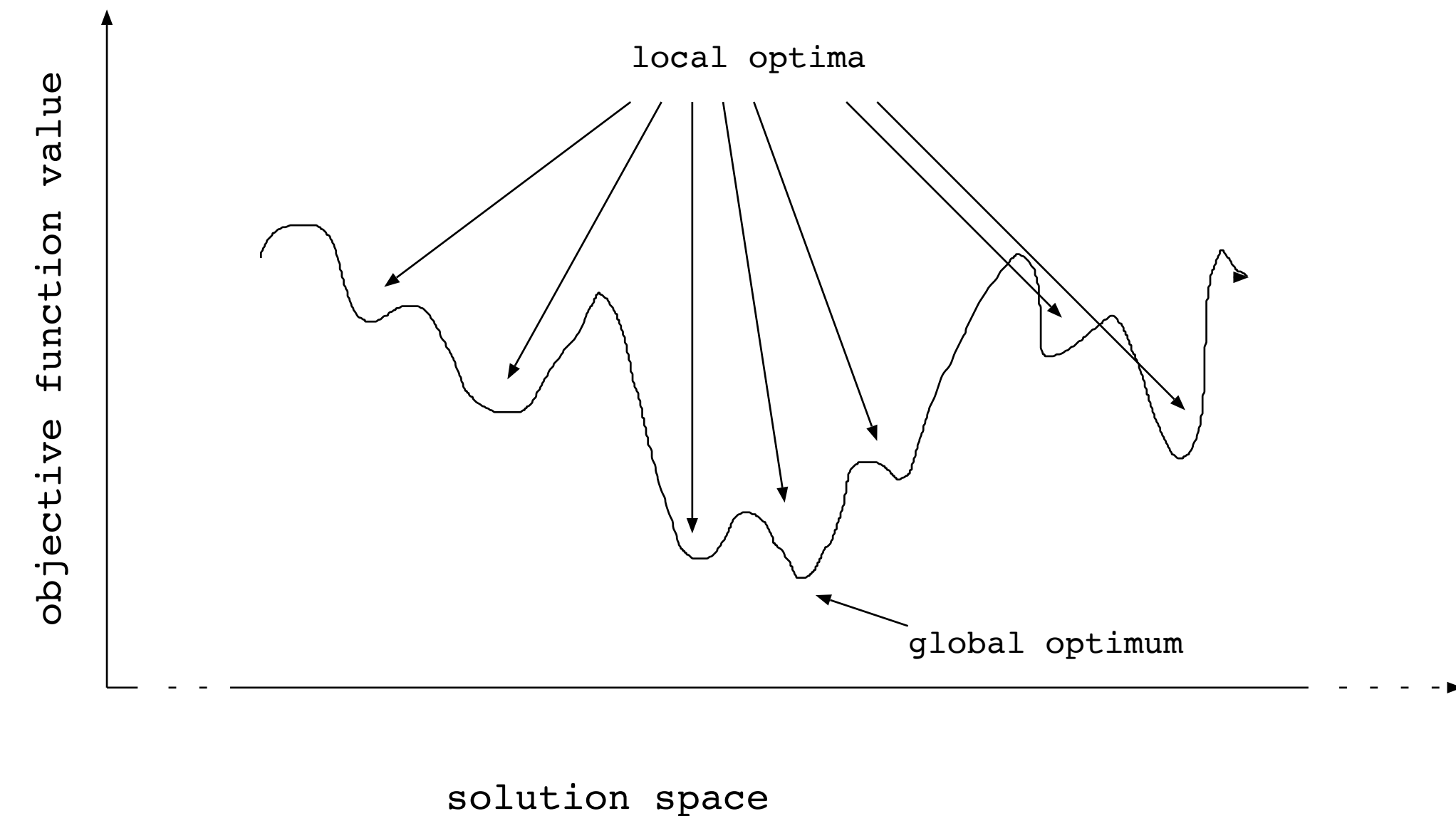
- neighbourhood size $\mathcal{O}(n^k)$



Problems with local search?

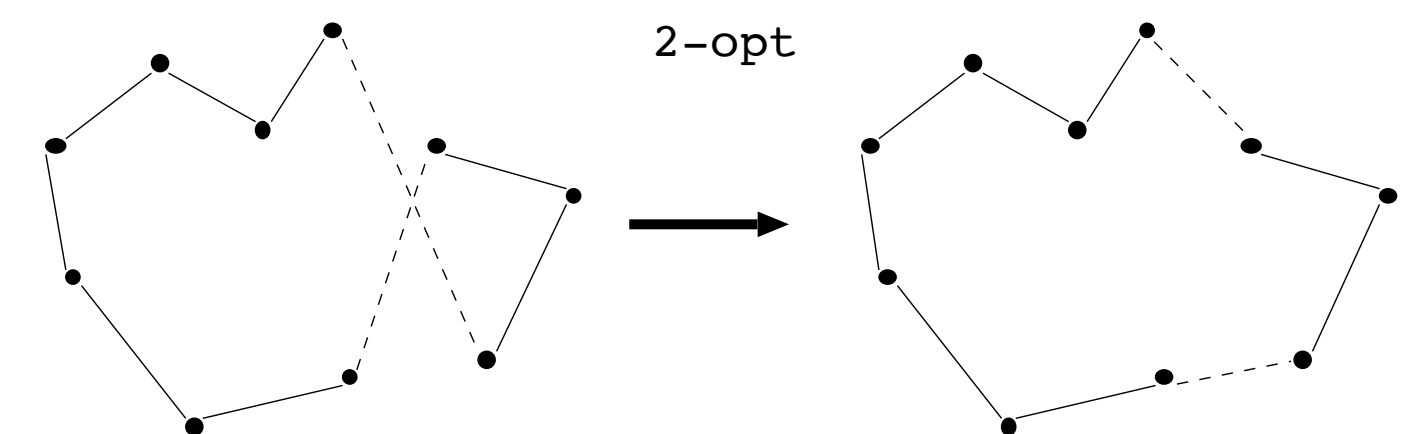
Stochastic Local Search:

- randomise initialisation step
 - random initial solutions
 - randomised construction heuristics
- randomise search steps
such that suboptimal/worsening steps are allowed
→ improved performance & robustness
- typically, degree of randomisation controlled by noise parameter
- allows to invest arbitrary computation times



Iterative Improvement for the TSP

- initial solution is a complete tour
- k -opt neighbourhood: solutions which differ by at most k edges



0-1 Knapsack Problem

There is a budget b available for investment in projects during the coming year and n projects are under consideration, where a_j is the outlay for project j and c_j is its expected return. The goal is to choose a set of projects so that the budget is not exceeded and the expected return is maximized.

Definition of the variables.

$x_j = 1$ if project j is selected, and $x_j = 0$ otherwise.

Definition of the constraints.

The budget cannot be exceeded:

$$\sum_{j=1}^n a_j x_j \leq b.$$

The variables are 0–1:

$$x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

Definition of the objective function.

The expected return is maximized:

$$\max \sum_{j=1}^n c_j x_j.$$

**Can you find a
feasible solution
greedily?**

Knapsack

Greedy algorithm

Definition of the objective function.

The expected return is maximized:

$$\max \sum_{j=1}^n c_j x_j.$$

The budget cannot be exceeded:

$$\sum_{j=1}^n a_j x_j \leq b.$$

The variables are 0–1:

$$x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

Sort items in increasing order of $\frac{c_j}{a_j}$

While budget $< b$ and there is an item that fits:

Insert next best item from sorted list into knapsack; update budget

Knapsack

Greedy algorithm

$$\max \sum_{j=1}^n c_j x_j. \quad \text{s.t.} \sum_{j=1}^n a_j x_j \leq b. \quad x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

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Insert next best item from sorted list into knapsack; update budget

$$b = 5$$

j	c_j	a_j	c_j/a_j
1	6	1	6
2	10	2	5
3	12	3	4

Knapsack

Greedy algorithm

$$\max \sum_{j=1}^n c_j x_j. \quad \text{s.t.} \sum_{j=1}^n a_j x_j \leq b. \quad x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

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$$b = 5$$

j	c_j	a_j	c_j/a_j	$c_j/a_j^{0.5}$
1	6	1	6	6
2	10	2	5	7.07
3	12	3	4	6.9

Branch & Bound for Integer Optimization

-LP-based $\min_x c^T x$ **s.t.** $Ax \leq b, x \in \{0,1\}^n$

Land & Doig, 1960

Repeat:

- 1 **Select Node**
- 2 **Solve LP Relaxation**
- 3 **Prune?**
- 4 **Add Cuts**
- 5 **Run Heuristics**
- 6 **Branch**

Branch & Bound for Integer Optimization

-LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

- 1 Select Node**
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- 6 Branch

N_0

Branch & Bound for Integer Optimization

-LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

$[0,1]^n$

Land & Doig, 1960

Repeat:

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N_0

→ Solve LP Relaxation
Lower Bound on OPT

Branch & Bound for Integer Optimization

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$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

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Land & Doig, 1960

Repeat:

- 1 **Select Node**
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- 4 Add Cuts
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N_0

Solve LP Relaxation
→ Lower Bound on OPT

worse than best solution?
Prune!

Branch & Bound for Integer Optimization

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Land & Doig, 1960

Repeat:

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N_0

Add Cuts:
Tightening Constraints

Branch & Bound for Integer Optimization

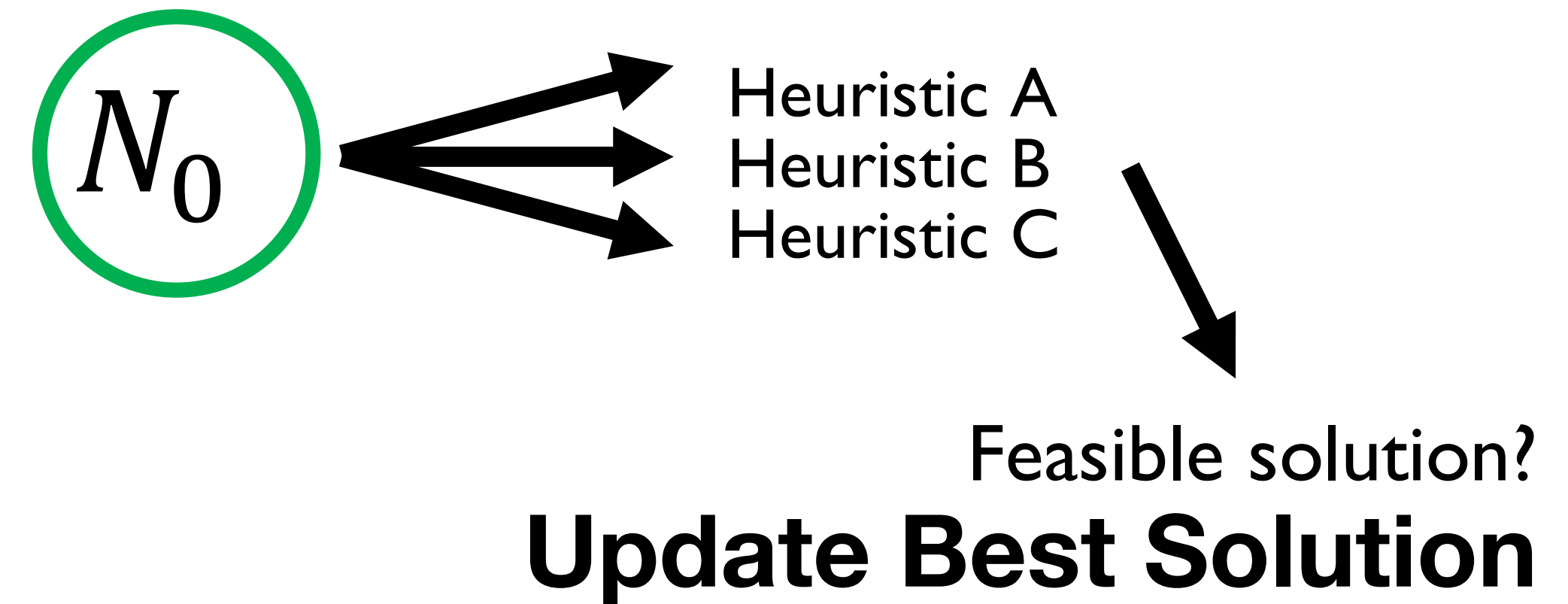
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Branch & Bound for Integer Optimization

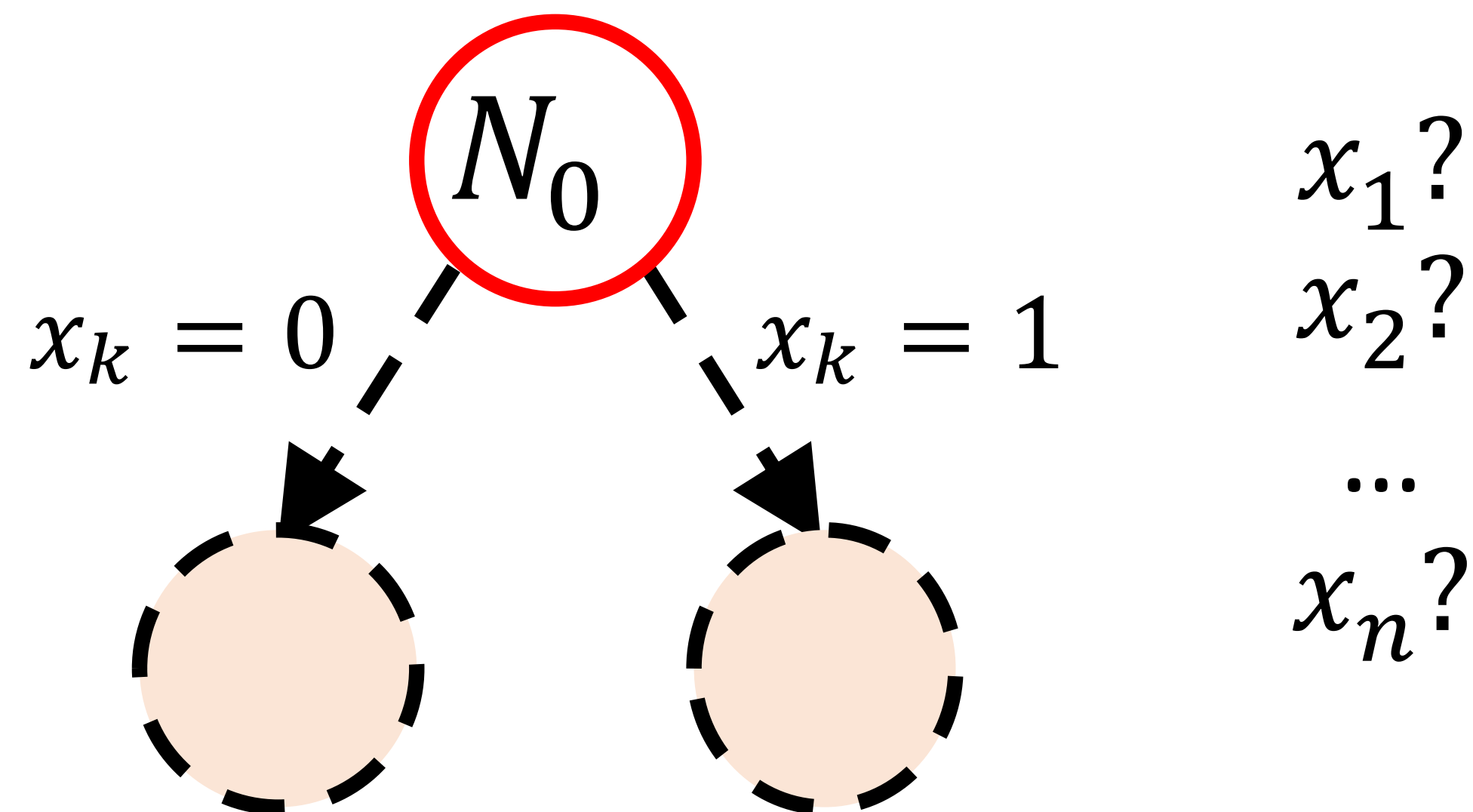
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Branch & Bound for Integer Optimization

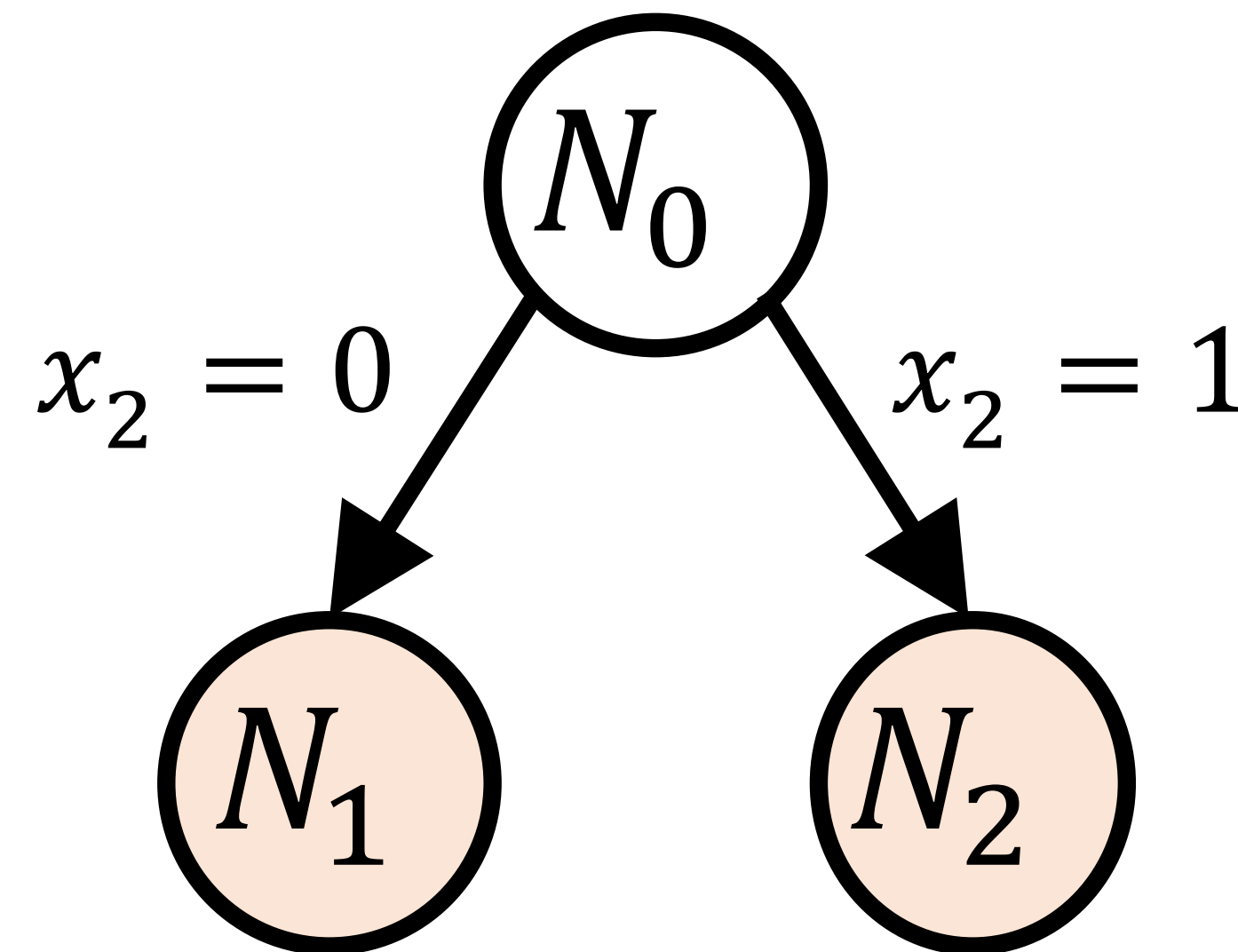
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Land & Doig, 1960

Repeat:

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$x_1?$
 x_2
 $\dots$
 $x_n?$

Branch & Bound for Integer Optimization

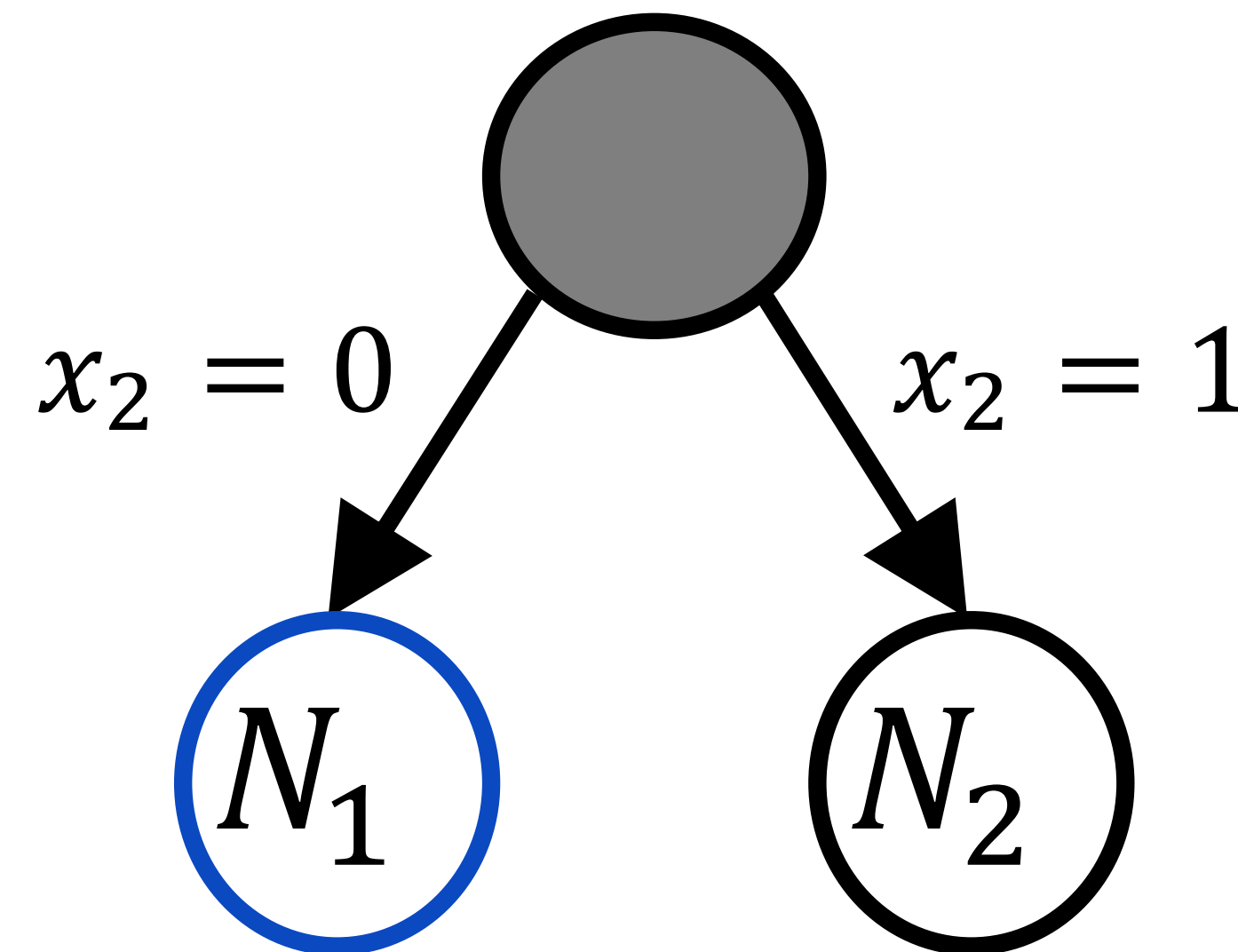
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Land & Doig, 1960

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Branch & Bound for Integer Optimization

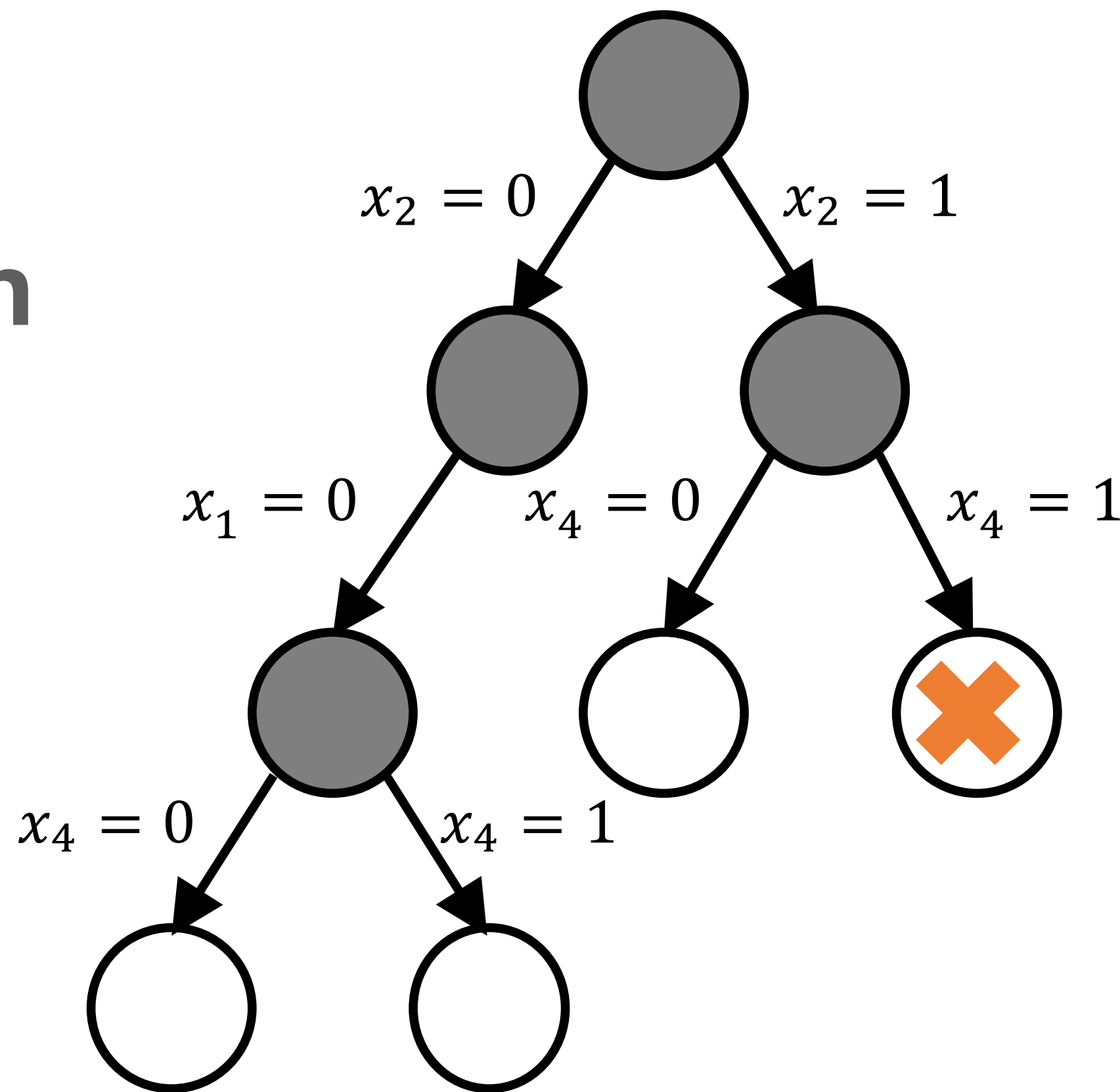
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Towards Tailored Algorithms



IBM Knowledge Center

Managing sets of parameters

Parameter names

Correspondence of parameters
between APIs

Saving parameter settings to a file
in the C API

– **Topical list of parameters**

Barrier

Benders algorithm

Distributed MIP

– **MIP**

MIP general

MIP strategies

MIP cuts

MIP tolerances

MIP limits

Here are links to parameters controlling MIP strategies.

[algorithm for initial MIP relaxation](#)

[Benders strategy](#)

[MIP subproblem algorithm](#)

[MIP variable selection strategy](#)

[MIP strategy best bound interval](#)

[MIP branching direction](#)

[backtracking tolerance](#)

[MIP dive strategy](#)

[MIP heuristic effort](#)

CPLEX Documentation

Towards Tailored Algorithms

MIP variable selection strategy

Value	Symbol	Meaning
-1	CPX_VARSEL_MININFEAS	Branch on variable with minimum infeasibility
0	CPX_VARSEL_DEFAULT	Automatic: let CPLEX choose variable to branch on; default
1	CPX_VARSEL_MAXINFEAS	Branch on variable with maximum infeasibility
2	CPX_VARSEL_PSEUDO	Branch based on pseudo costs
3	CPX_VARSEL_STRONG	Strong branching
4	CPX_VARSEL_PSEUDOREDUCED	Branch based on pseudo reduced costs

MIP heuristic frequency

Value	Meaning
-1	None
0	Automatic: let CPLEX choose; default
Any positive integer	Apply the periodic heuristic at this frequency

- an algorithm A with parameters $p_1, \dots, p_k$ that affect its behaviour,
- a space C of parameter settings (configurations), where $c \in C$ specifies values for $p_1, \dots, p_k$,
- a set of problem instances I ,
- a performance metric m that measures the performance of A on instance set I for a given configuration c ,

find a configuration $c^* \in C$ such that running algorithm A on instance set I maximizes metric m

Greedy 0-1 knapsack Exponent p_1 such that items are sorted w.r.t. $c_j/a_j^{p_1}$

- an algorithm A with parameters $p_1, \dots, p_k$ that affect its behaviour,
- a space C of parameter settings (configurations), where $c \in C$ specifies values for $p_1, \dots, p_k$,
 $p_1 \in (0,1]$
- a set of problem instances I ,
- a performance metric m that measures the performance of A on instance set I for a given configuration c ,

$$\sum_{i \in I} \sum_{j=1}^n c_j x_j$$

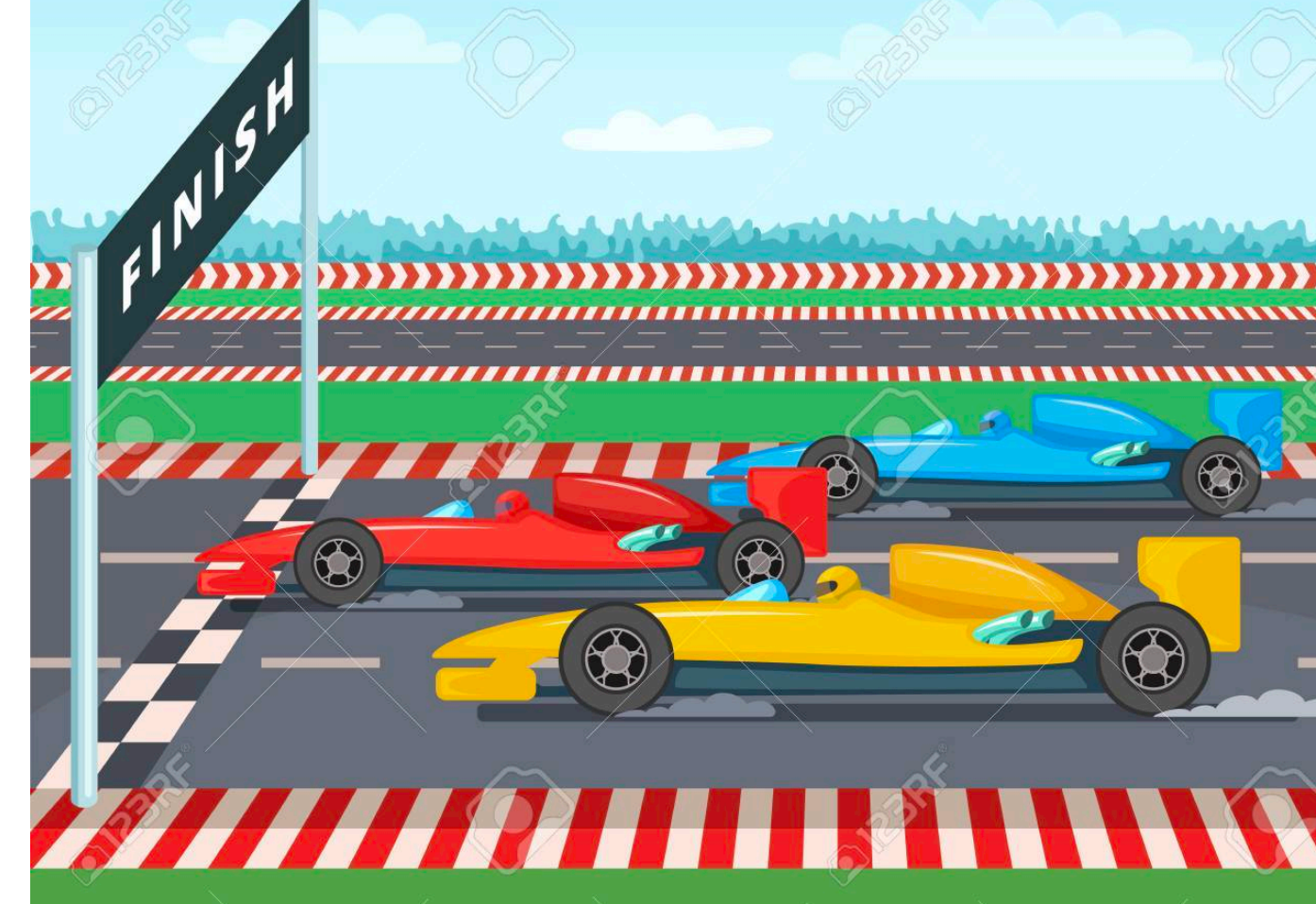
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Issues to consider

- Generalization
- Time-outs!
- **Optimization!**

Racing Procedures

- Assumptions: small, finite configuration space C
- Basic idea:
 - sample instance,
 - test remaining configs.,
 - eliminate really bad ones relative to current best config.,
 - repeat.



123rf.com

procedure *F-Race*

input *target algorithm A, set of configurations C, set of problem instances I,*
performance metric m;

parameters *integer ni_{min} ;*

output *set of configurations C^* ;*

$C^* := C$; $ni := 0$;

repeat

 randomly choose instance i from set I ;

 run all configurations of A in C^* on i ;

$ni := ni + 1$;

if $ni \geq ni_{min}$ **then**

 perform rank-based Friedman test on results for configurations in C^* on all instances
 in I evaluated so far;

if test indicates significant performance differences **then**

$c^* :=$ best configuration in C^* (according to m over instances evaluated so far);

for all $c \in C^* \setminus \{c^*\}$ **do**

 perform pairwise Friedman post hoc test on c and c^* ;

if test indicates significant performance differences **then**

 eliminate c from C^* ;

end if;

end for;

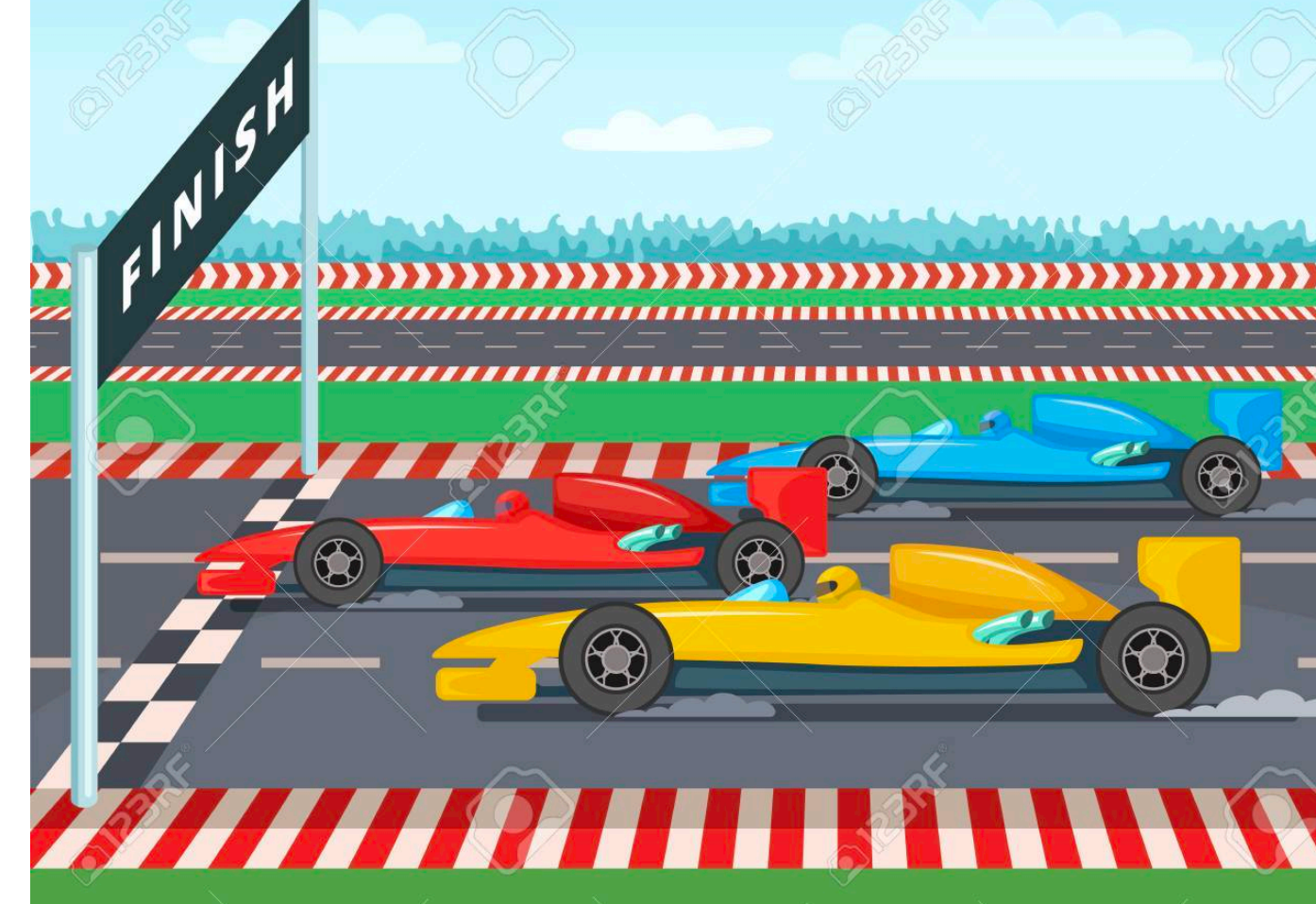
end if;

end if;

until termination condition met;

return C^* ;

end *F-Race*



123rf.com

Birattari, Mauro, et al. "A Racing Algorithm
for Configuring Metaheuristics." GECCO.
Vol. 2. No. 2002.

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end if;

end for;

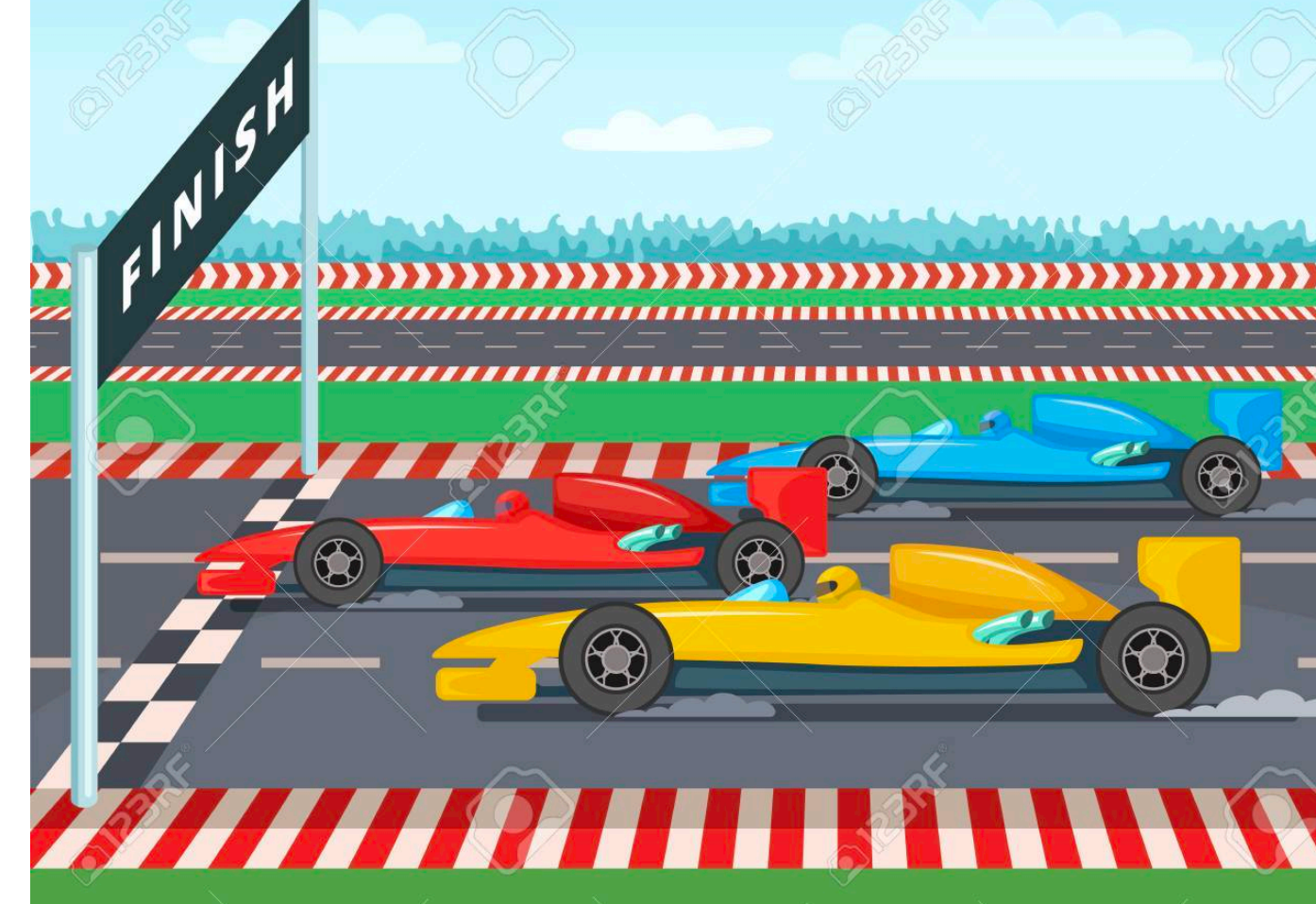
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123rf.com

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for Configuring Metaheuristics." GECCO.
Vol. 2. No. 2002.

procedure *F-Race*

input target algorithm A , set of configurations C , set of problem instances I ,
performance metric m ;

parameters integer ni_{min} ;

output set of configurations C^* ;

$C^* := C$; $ni := 0$;

repeat

randomly choose instance i from set I ;

run all configurations of A in C^* on i ;

$ni := ni + 1$;

if $ni \geq ni_{min}$ **then**

perform rank-based Friedman test on results for configurations in C^* on all instances
in I evaluated so far;

if test indicates significant performance differences **then**

$c^* :=$ best configuration in C^* (according to m over instances evaluated so far);

for all $c \in C^* \setminus \{c^*\}$ **do**

perform pairwise Friedman post hoc test on c and c^* ;

if test indicates significant performance differences **then**

eliminate c from C^* ;

end if;

end for;

end if;

end if;

until termination condition met;

return C^* ;

end *F-Race*

Are these results
statistically “similar”?

	Config X	Config Y	Config Z
Instance 1	7	5	8
Instance 2	1	4	2
Instance 3	3	3	3
Instance 4	2	7	8

Sum of ranks	4	7	9
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$$T = \frac{(n-1) \sum_{j=1}^n \left(R_j - \frac{k(n+1)}{2} \right)^2}{\sum_{l=1}^k \sum_{j=1}^n R_{lj}^2 - \frac{kn(n+1)^2}{4}}$$

Birattari, Mauro, et al. "A Racing Algorithm
for Configuring Metaheuristics." GECCO.
Vol. 2. No. 2002.

procedure *I/F-Race*

input *target algorithm A, set of configurations C, set of problem instances I,*
performance metric m;

output *set of configurations C*;*

initialise probabilistic model M;

C' := ∅; // later, C' is the set of survivors from the previous F-Race

repeat

based on model M, sample set of configurations $\hat{C} \subseteq C$;

perform F-Race on configurations in $\hat{C} \cup C'$ to obtain set of configurations C;*

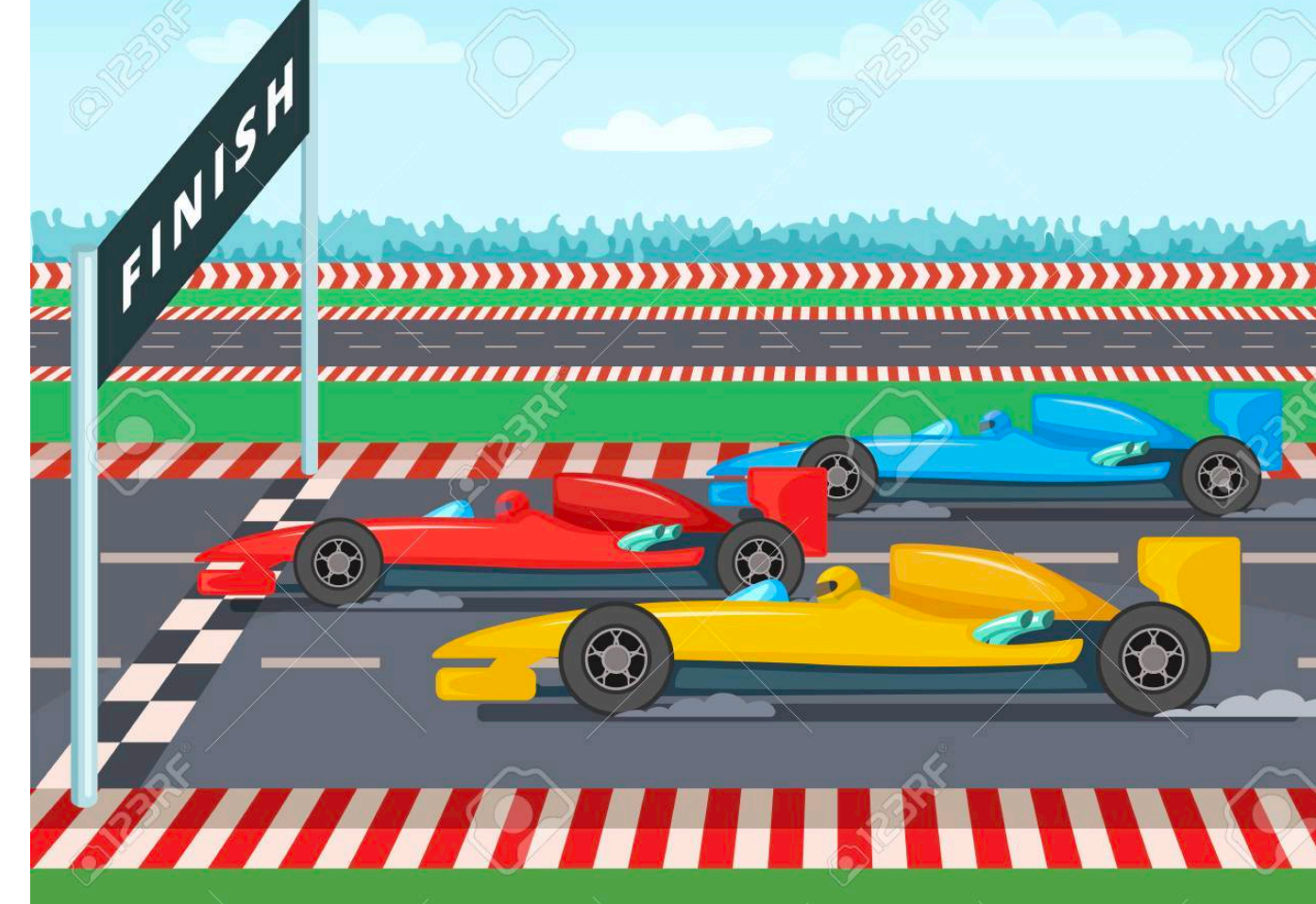
update probabilistic model M based on configurations in C;*

C' := C;*

until *termination condition met;*

return *c* ∈ C* with best performance (according to m) over all instances evaluated;*

end *I/F-Race*



123rf.com

Balaprakash, Prasanna, Mauro Birattari, and Thomas Stützle. "Improvement strategies for the F-Race algorithm: Sampling design and iterative refinement." International workshop on hybrid metaheuristics. Springer, Berlin, Heidelberg, 2007.

ParamILS

procedure *ParamILS*

input *target algorithm A, set of configurations C, set of problem instances I,*
performance metric m;

parameters *configuration $c_0 \in C$, integer r, integer s, probability pr;*

output *configuration c^* ;*

$c^* := c_0$;

for $i := 1$ **to** r **do**

draw c from C uniformly at random;

assess c against c^* based on performance of A on instances from I according to metric m ;

if c found to perform better than c^* **then**

$c^* := c$;

end if;

end for;

$c := c^*$;

perform subsidiary local search on c ;

while termination condition not met **do**

$c' := c$;

perform s random perturbation steps on c'

perform subsidiary local search on c' ;

assess c' against c based on performance of A on instances from I according to metric m ;

if c' found to perform better than c **then** *// acceptance criterion*

update overall incumbent c^* ;

$c := c'$;

end if;

with probability pr **do**

draw c from C uniformly at random;

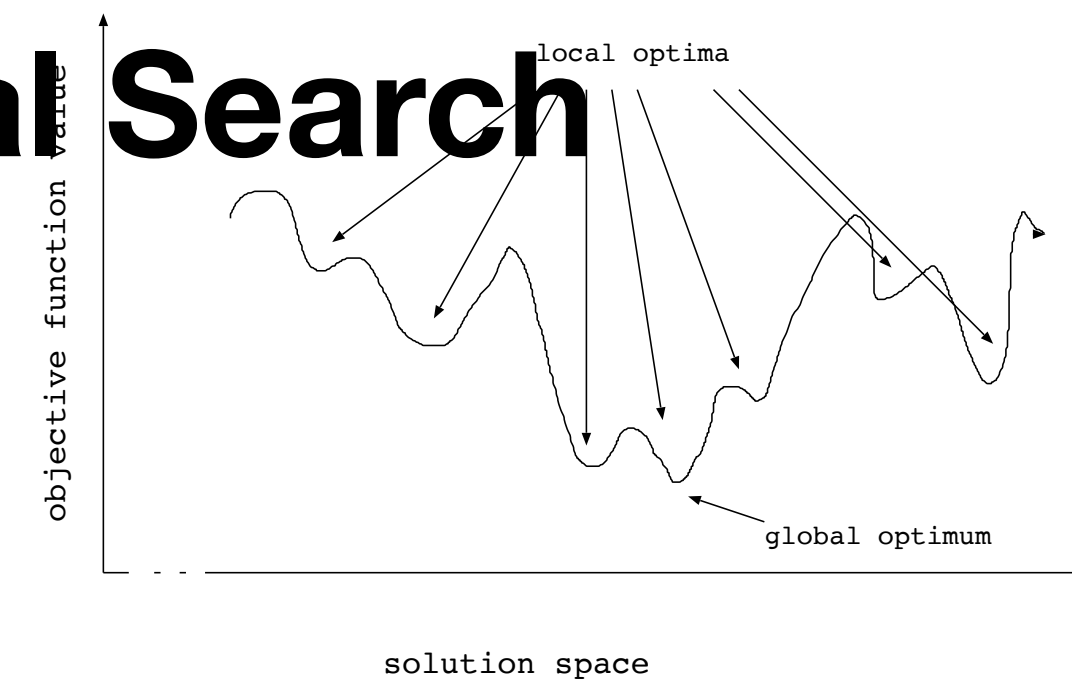
end with probability;

end while;

return c^* ;

end *ParamILS*

ILS: Iterated Local Search



Initial sampling phase

Random perturbation + local search
Evaluation

Update incumbent config.

Random restart!

Hutter, Frank, et al. "ParamILS: an automatic algorithm configuration framework." *Journal of Artificial Intelligence Research* 36 (2009): 267-306.

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$c' := c$;

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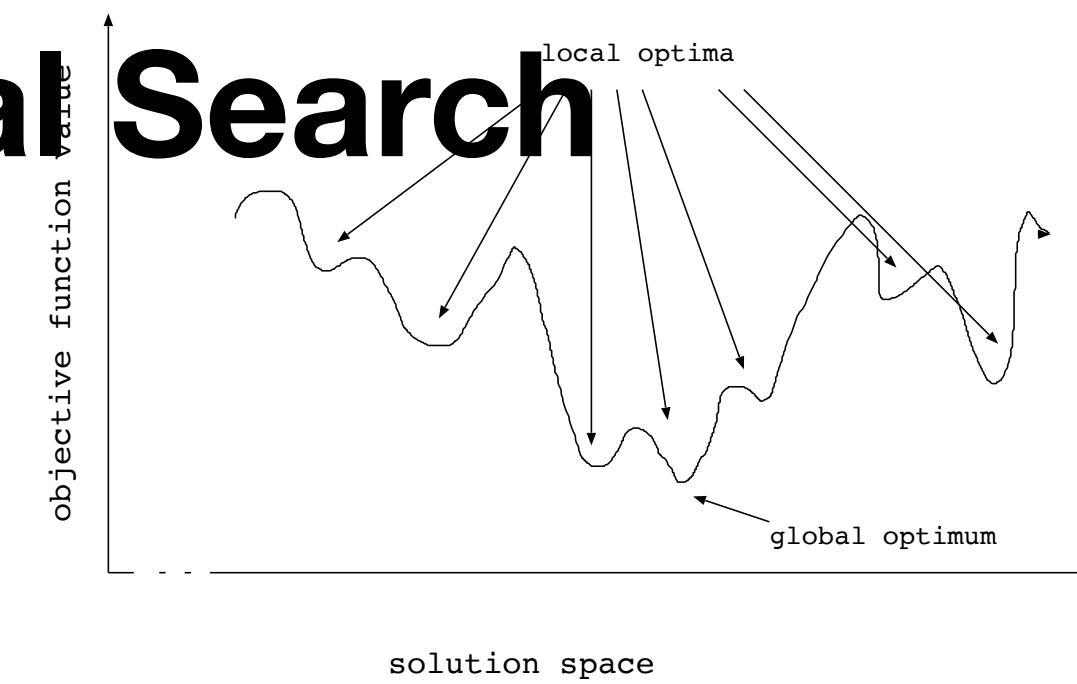
end with probability;

end while;

return c^* ;

end *ParamILS*

ILS: Iterated Local Search



Initial sampling phase

Random perturbation + local search
Evaluation

Update incumbent config.

Random restart!

Hutter, Frank, et al. "ParamILS: an automatic algorithm configuration framework." *Journal of Artificial Intelligence Research* 36 (2009): 267-306.

Sequential Model-Based Optimization

procedure *SMBO*

input *target algorithm A, set of configurations C, set of problem instances I,*
performance metric m;

output *configuration c^* ;*

determine initial set of configurations $C_0 \subset C$;

for all $c \in C_0$, measure performance of *A* on *I* according to metric *m*;

build initial model *M* based on performance measurements for C_0 ;

determine incumbent $c^* \in C_0$ for which best performance was observed or predicted;

repeat

based on model *M*, determine set of configurations $C' \subseteq C$;

for all $c \in C'$, measure performance of *A* on *I* according to metric *m*;

update model *M* based on performance measurements for C' ;

update incumbent c^* ;

until termination condition met;

return c^* ;

end *SMBO*

Sequential Model-Based Optimization

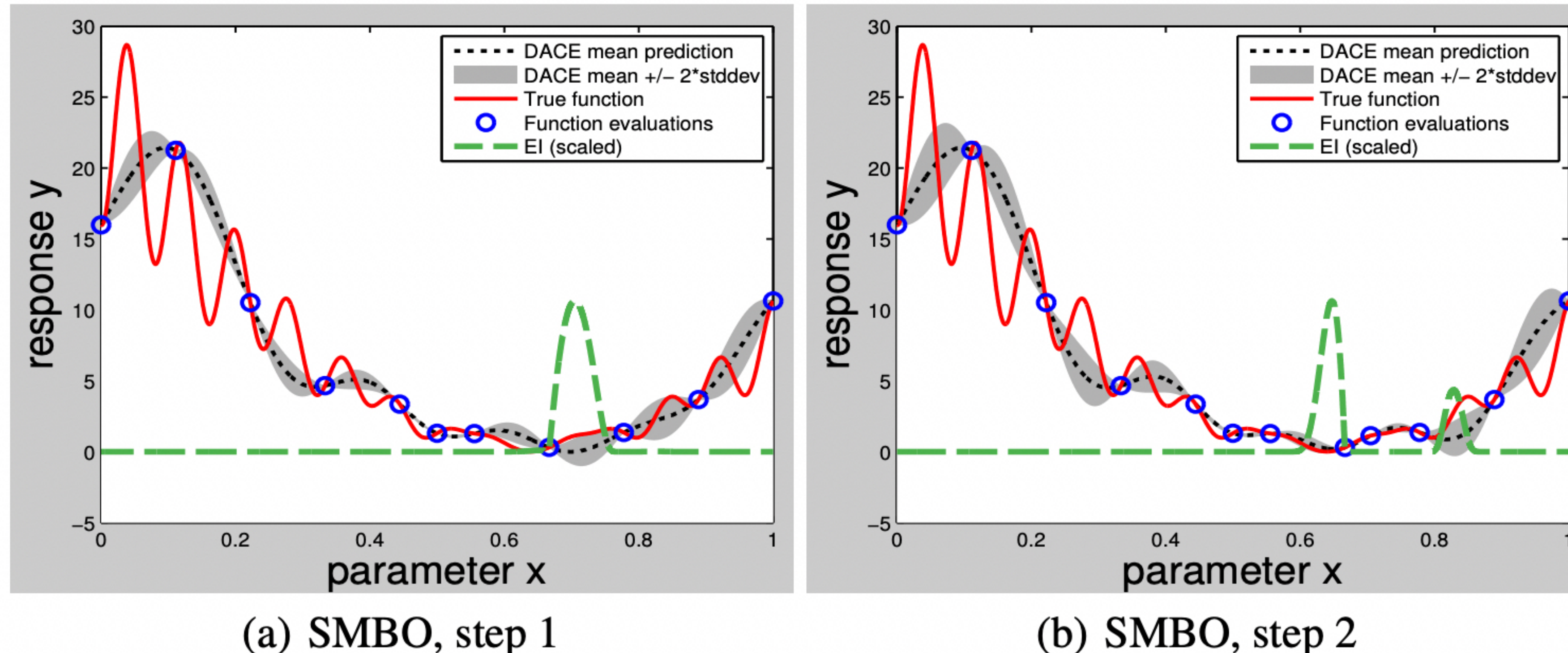


Fig. 1. Two steps of SMBO for the optimization of a 1D function. The true function is shown as a solid line, and the circles denote our observations. The dotted line denotes the mean prediction of a noise-free Gaussian process model (the “DACE” model), with the grey area denoting its uncertainty. Expected improvement (scaled for visualization) is shown as a dashed line.

Sequential Model-Based Optimization

procedure *SMBO*

input *target algorithm A, set of configurations C, set of problem instances I,*
performance metric m;

output *configuration c^* ;*

determine initial set of configurations $C_0 \subset C$;

for all $c \in C_0$, measure performance of *A* on *I* according to metric *m*;

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determine incumbent $c^* \in C_0$ for which best performance was observed or predicted;

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based on model *M*, determine set of configurations $C' \subseteq C$;

for all $c \in C'$, measure performance of *A* on *I* according to metric *m*;

update model *M* based on performance measurements for C' ;

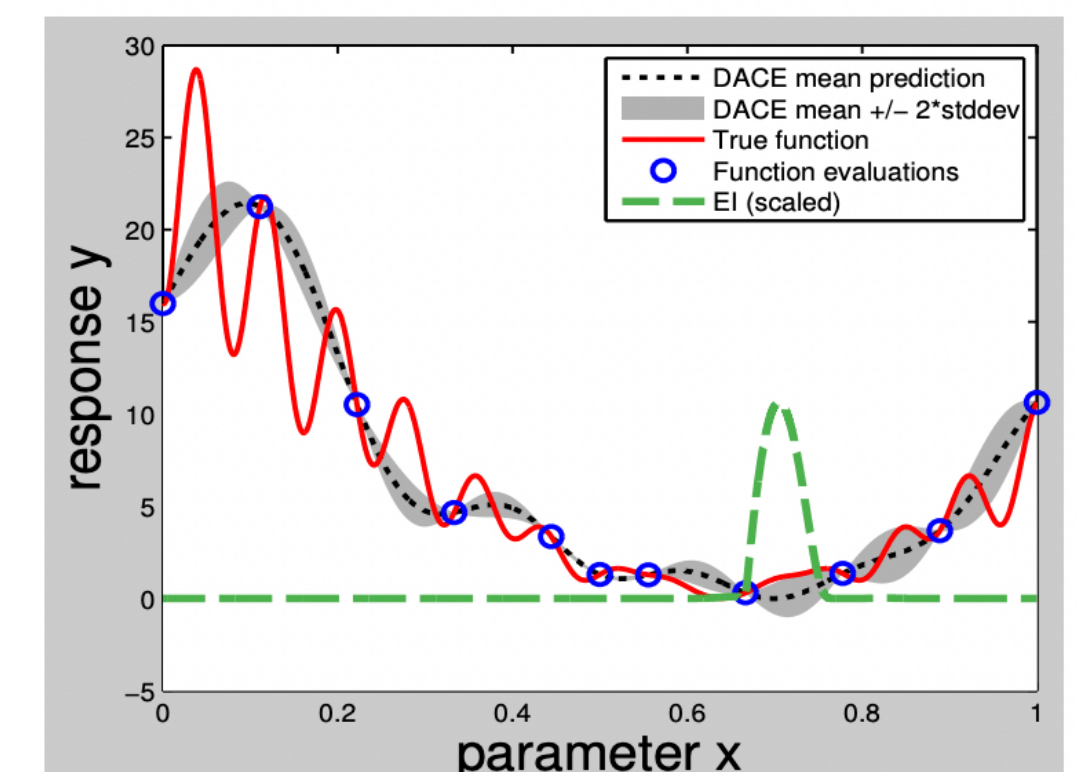
update incumbent c^* ;

until termination condition met;

return c^* ;

end *SMBO*

$$EI(x) := E[\max\{f_{min} - \hat{F}(x), 0\}]$$



Other flavours of algo. config.

“

- **per-instance algorithm selection** methods choose one of several target algorithms to be applied to a given problem instance based on properties of that instance determined just before attempting to solve it
- **Reactive search procedures**, on-line algorithm control methods and adaptive operator selection techniques switch between different algorithms, heuristic mechanisms and parameter configurations while running on a given problem instance
- **dynamic algorithm portfolio** approaches repeatedly adjust the allocation of CPU shares between algorithms that are running concurrently on a given problem instance

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